

Generalized Hermite Reduction, Creative Telescoping, and Definite Integration of D-Finite Functions

Frédéric Chyzak



February 10–14, 2020, DART X (New York City, USA)
The 10th international workshop on Differential Algebra and Related Topics

Based on joint work with A. Bostan, P. Lairez, and B. Salvy

Parametrized Definite Integrals

$$\int_0^{+\infty} x J_1(ax) I_1(ax) Y_0(x) K_0(x) dx = -\frac{\ln(1-a^4)}{2\pi a^2} \quad (\text{Glasser \& Montaldi, 1994})$$

$$\int_0^\infty \int_0^\infty J_1(x) J_1(y) J_2(c\sqrt{xy}) \frac{dx dy}{e^{x+y}} \quad (\text{has a 2nd-order linear ODE})$$

$$\frac{1}{2\pi i} \oint \frac{(1+2xy+4y^2) \exp\left(\frac{4x^2y^2}{1+4y^2}\right)}{y^{n+1}(1+4y^2)^{\frac{3}{2}}} dy = \frac{H_n(x)}{[n/2]!} \quad (\text{Doetsch, 1930})$$

$$\int_{-1}^{+1} \frac{e^{-px} T_n(x)}{\sqrt{1-x^2}} dx = (-1)^n \pi I_n(p)$$

$$\int_0^{+\infty} x e^{-px^2} J_n(bx) I_n(cx) dx = \frac{1}{2p} \exp\left(\frac{c^2 - b^2}{4p}\right) J_n\left(\frac{bc}{2p}\right)$$

$$\frac{1}{(2i\pi)^2} \oint \oint \frac{f(s, t/s, x/t)}{st} ds dt = 1 + 6 \cdot \int_0^x \frac{{}_2F_1\left(\begin{matrix} 1/3, 2/3 \\ 2 \end{matrix} \middle| \frac{27w(2-3w)}{(1-4w)^3}\right)}{(1-4w)(1-64w)} dw$$

$$\text{where } f(s, t, u) = \frac{(1-s)(1-t)(1-u)}{1-2(s+t+u)+3(st+tu+us)-4stu}$$

(Bostan, Chyzak, van Hoeij, Pech, 2011)

Differentiating under the Integral Sign

Zeilberger's derivation (1982) of a classical integral

Given $f(b, x) = e^{-x^2} \cos 2bx$, find $F(b) = \int_{-\infty}^{+\infty} f(b, x) dx = ?$.

$$\frac{dF}{db}(b) = \int_{-\infty}^{+\infty} -2xe^{-x^2} \sin 2bx dx =$$

$$\left[e^{-x^2} \sin 2bx \right]_{x=-\infty}^{x=+\infty} + \int_{-\infty}^{+\infty} -2be^{-x^2} \cos 2bx dx = -2b F(b).$$

Differentiating under the Integral Sign

Zeilberger's derivation (1982) of a classical integral

Given $f(b, x) = e^{-x^2} \cos 2bx$, find $F(b) = \int_{-\infty}^{+\infty} f(b, x) dx = ?$.

$$\begin{aligned} \frac{dF}{db}(b) &= \int_{-\infty}^{+\infty} -2xe^{-x^2} \sin 2bx dx = \\ &= \left[e^{-x^2} \sin 2bx \right]_{x=-\infty}^{x=+\infty} + \int_{-\infty}^{+\infty} -2be^{-x^2} \cos 2bx dx = -2bF(b). \end{aligned}$$

Continuous form of “Creative Telescoping”:

$$\frac{df}{db}(b, x) + 2bf(b, x) = \frac{dg}{dx}(b, x) \quad \text{for} \quad g(b, x) = -\frac{1}{2x} \frac{df}{db}(b, x).$$

After integration over x from $-\infty$ to $+\infty$:

$$\frac{dF}{db}(b) + 2bF(b) = \left[\frac{dg}{dx}(b, x) \right]_{x=-\infty}^{x=+\infty} = 0 - 0 = 0.$$

Hermite Reduction (1872)

$$EA - mFA' = P \implies \int \frac{P}{A^{m+1}} = \frac{F}{A^m} + \int \frac{E + F'}{A^m}$$

Cela posé, l'intégrale $\int \frac{P dx}{A^{m+1}}$ se traitera comme il suit : nous effectuerons sur A et sa dérivée A' les opérations du plus grand commun diviseur, de manière à obtenir deux polynômes G et H , satisfaisant à la condition

$$AG - A'H = 1.$$

Nous formerons ensuite deux séries de fonctions entières :

$$V_0, V_1, \dots, V_{m-1},$$

$$P_1, P_2, \dots, P_m,$$

par ces relations, où les polynômes Q_1, Q_2, \dots sont entièrement arbitraires, savoir :

$$\begin{aligned} \alpha V_1 &= HP - AQ_1, \\ (\alpha - 1)V_2 &= HP - AQ_2, \\ (\alpha - 2)V_3 &= HP - AQ_3, \\ &\dots\dots\dots \\ V_{m-1} &= HP_{m-1} - AQ_{m-1}, \\ P_1 &= GP - A'Q - V_1, \\ P_2 &= GP - A'Q_2 - V_2, \\ &\dots\dots\dots \\ P_m &= GP_{m-1} - A'Q_{m-1} - V_{m-1}. \end{aligned}$$

Maintenant je prouverai qu'en faisant

$$\begin{aligned} V &= V_0 + AV_1 + A^2V_2 + \dots + A^{m-1}V_{m-1}, \\ U &= P_m, \end{aligned}$$

on a l'égalité

$$\frac{P}{A^{m+1}} = \frac{U}{A} + \left(\frac{V}{A^2}\right)'$$

d'où

$$\int \frac{P dx}{A^{m+1}} = \int \frac{U dx}{A} + \frac{V}{A^2},$$

de sorte que $\frac{V}{A^2}$ est la partie algébrique de l'intégrale proposée, et $\int \frac{U dx}{A}$ la partie transcendante.

A cet effet, j'élimine G et H entre les trois égalités

$$\begin{aligned} AG - A'H &= 1, \\ (\alpha - i)V_i &= HP_i - AQ_i, \\ P_{i+1} &= GP_i - A'Q_i - V_i, \end{aligned}$$

ce qui donne

$$AP_{i+1} = P_i + (\alpha - i)A'V_i - AV_i.$$

On en peut écrire cette relation de la manière suivante :

$$\frac{P_i}{A^{m-i+1}} - \frac{P_{i+1}}{A^{m-i}} = \left(\frac{V_i}{A^{m-i}}\right)'$$

En supposant ensuite $i = 0, 1, 2, \dots, \alpha - 1$ et ajoutant membre à membre, nous en concluons

$$\frac{P}{A^{m+1}} - \frac{P_\alpha}{A} = \left(\frac{V_0}{A^2} + \frac{V_1}{A^{m-1}} + \dots + \frac{V_{\alpha-1}}{A}\right)'$$

ce qui fait bien voir qu'on satisfait à la condition proposée

$$\frac{P}{A^{m+1}} = \frac{U}{A} + \left(\frac{V}{A^2}\right)'$$

par les valeurs

$$\begin{aligned} V &= V_0 + AV_1 + A^2V_2 + \dots + A^{m-1}V_{m-1}, \\ U &= P_m, \end{aligned}$$

comme il s'agissait de le démontrer.

Ostrogradsky–Hermite Reduction

$$EA - mFA' = P \implies \int \frac{P}{A^{m+1}} = \frac{F}{A^m} + \int \frac{E + F'}{A^m}$$

Cela posé, l'intégrale $\int \frac{P dx}{A^{m+1}}$ se traitera comme il suit : nous effectuerons sur A et sa dérivée A' les opérations du plus grand commun diviseur, de manière à obtenir deux polynômes G et H , satisfaisant à la condition

$$AG - A'H = 1.$$

Nous ferons ensuite deux séries de fonctions entières :

$$\begin{aligned} V_0, V_1, \dots, V_{m-1}, \\ P_1, P_2, \dots, P_m. \end{aligned}$$

par ces relations, où les polynômes Q_1, Q_2, \dots sont entièrement arbitraires, savoir :

$$\begin{aligned} \alpha V_1 &= HP - AQ_1, \\ (\alpha - 1)V_1 &= HP - AQ_1, \\ (\alpha - 2)V_1 &= HP - AQ_1, \\ &\dots\dots\dots \\ V_{m-1} &= HP_{m-1} - AQ_{m-1}, \\ P_1 &= GP - A'Q - V_1, \\ P_2 &= GP_1 - A'Q_1 - V_1, \\ &\dots\dots\dots \\ P_m &= GP_{m-1} - A'Q_{m-1} - V_{m-1}. \end{aligned}$$

Maintenant je prouverai qu'en faisant

$$\begin{aligned} V &= V_0 + AV_1 + A^2V_2 + \dots + A^{m-1}V_{m-1}, \\ U &= P_m, \end{aligned}$$

on a l'égalité

$$\frac{P}{A^{m+1}} = \frac{U}{A} + \left(\frac{V}{A^2}\right)'$$

d'où

$$\int \frac{P dx}{A^{m+1}} = \int \frac{U dx}{A} + \frac{V}{A^2},$$

de sorte que $\frac{V}{A^2}$ est la partie algébrique de l'intégrale proposée, et $\int \frac{U dx}{A}$ la partie transcendante.

A cet effet, j'élimine G et H entre les trois égalités

$$\begin{aligned} AG - A'H &= 1, \\ (\alpha - i)V_i &= HP_i - AQ_i, \\ P_{i+1} &= GP_i - A'Q_i - V_i, \end{aligned}$$

ce qui donne

$$AP_{i+1} = P_i + (\alpha - i)A'V_i - AV_i.$$

On en peut écrire cette relation de la manière suivante :

$$\frac{P_i}{A^{m-i+1}} - \frac{P_{i+1}}{A^{m-i}} = \left(\frac{V_i}{A^{m-i}}\right)'$$

En supposant ensuite $i = 0, 1, 2, \dots, \alpha - 1$ et ajoutant membre à membre, nous en concluons

$$\frac{P}{A^{m+1}} - \frac{P_\alpha}{A} = \left(\frac{V_0}{A^2} + \frac{V_1}{A^{m-1}} + \dots + \frac{V_{\alpha-1}}{A}\right)'$$

ce qui fait bien voir qu'on satisfait à la condition proposée

$$\frac{P}{A^{m+1}} = \frac{U}{A} + \left(\frac{V}{A^2}\right)'$$

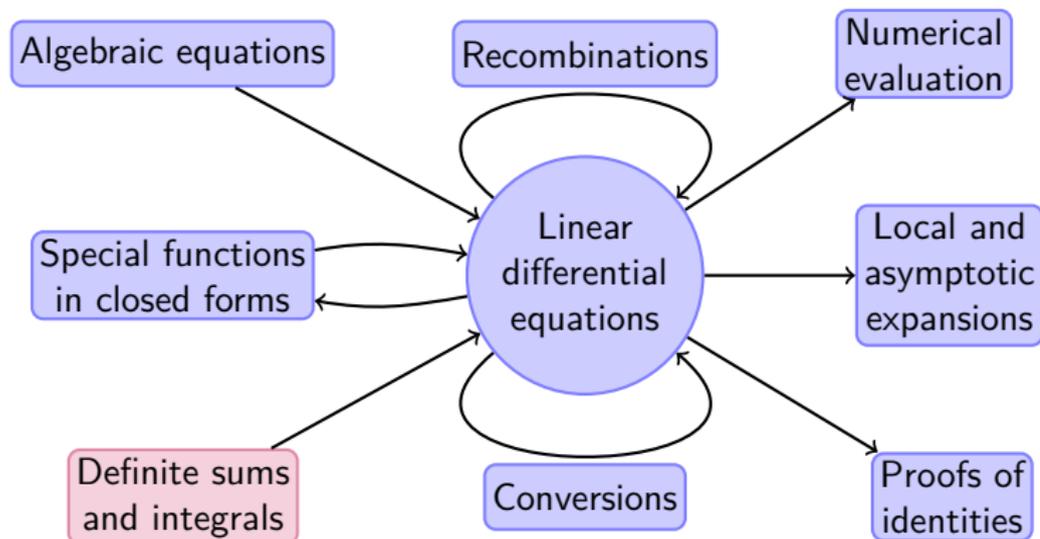
par les valeurs

$$\begin{aligned} V &= V_0 + AV_1 + A^2V_2 + \dots + A^{m-1}V_{m-1}, \\ U &= P_m, \end{aligned}$$

comme il s'agissait de le démontrer.

See also (Ostrogradsky, 1833, 1844/45).

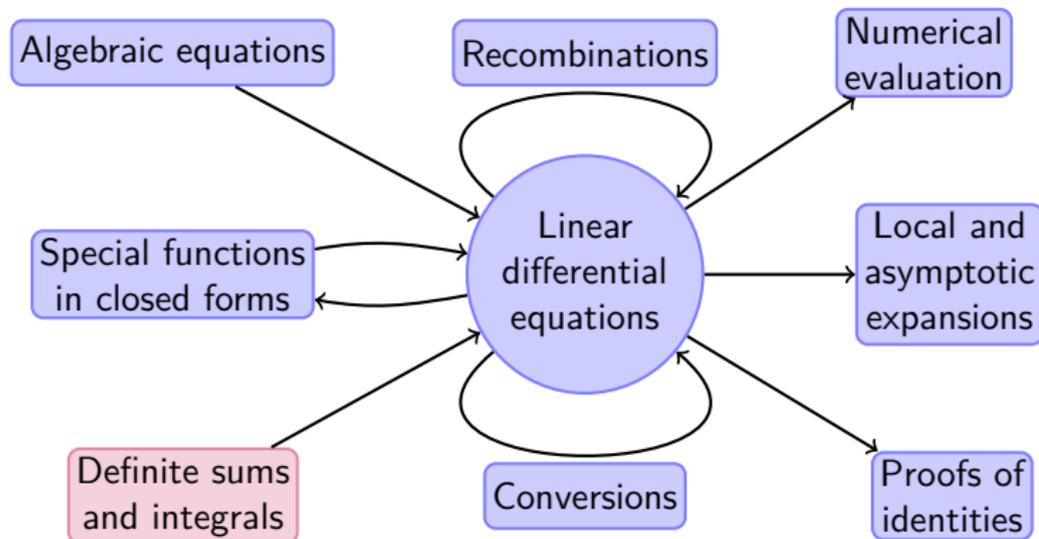
Linear Differential Equations as a Data Structure



Def: differentially finite functions (a.k.a. **D-finite**)

A function $f(x)$ is D-finite if its derivatives $f(x), f'(x), f''(x), \dots$, span a finite-dimensional vector space over $\mathbb{C}(x)$.

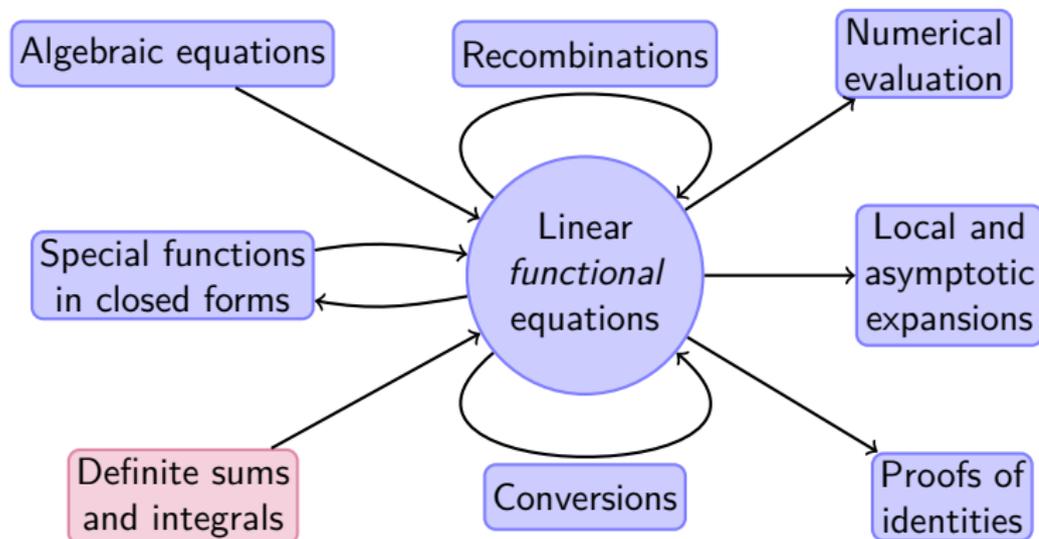
Linear Differential Equations as a Data Structure



Def: multivariate D-finite functions

A function $f(x, y, z)$ is D-finite iff its derivatives $\frac{\partial^{i+j+k} f}{\partial x^i \partial y^j \partial z^k}(x, y, z)$, $i, j, k \geq 0$, span a finite-dimensional vector space over $\mathbb{C}(x, y, z)$.

Linear Differential Equations as a Data Structure



Def: multivariate ∂ -finite functions

A function $f_{n,m}(x, y, z)$ is ∂ -finite iff a similar confinement holds for derivatives w.r.t. x, y, z , shifts w.r.t. n, m , etc.

Creative Telescoping for Sums and Integrals

$$U_n = \sum_{k=a}^b u_{n,k} = ?$$

Given a relation $a_r(n)u_{n+r,k} + \dots + a_0(n)u_{n,k} = v_{n,k+1} - v_{n,k}$, summation leads by “telescoping” to

$$a_r(n)U_{n+r} + \dots + a_0(n)U_n = v_{n,b+1} - v_{n,a} \stackrel{\text{often}}{=} 0.$$

$$U(t) = \int_a^b u(t, x) dx = ?$$

Given a relation $a_r(t)\frac{\partial^r u}{\partial t^r} + \dots + a_0(t)u = \frac{\partial}{\partial x}v(t, x)$, integrating leads by “telescoping” to

$$a_r(t)\frac{\partial^r U}{\partial t^r} + \dots + a_0(t)U = v(t, b) - v(t, a) \stackrel{\text{often}}{=} 0.$$

Adapts easily to $U(t) = \sum_{k=a}^b u_k(t)$, $U_n = \int_a^b u_n(x) dx$, etc.

Creative Telescoping for Sums and Integrals

$$U_n = \sum_{k=a}^b u_{n,k} = ?$$

Given a relation $a_r(n)u_{n+r,k} + \dots + a_0(n)u_{n,k} = v_{n,k+1} - v_{n,k}$, summation leads by “telescoping” to

$$a_r(n)U_{n+r} + \dots + a_0(n)U_n = v_{n,b+1} - v_{n,a} \stackrel{\text{often}}{=} 0.$$

$$U(t) = \int_a^b u(t,x) dx = ?$$

Given a relation $a_r(t)\frac{\partial^r u}{\partial t^r} + \dots + a_0(t)u = \frac{\partial}{\partial x}v(t,x)$, integrating leads by “telescoping” to

$$a_r(t)\frac{\partial^r U}{\partial t^r} + \dots + a_0(t)U = v(t,b) - v(t,a) \stackrel{\text{often}}{=} 0.$$

Telescopier

Certificate

History of Algorithms for Creative Telescoping

Algorithmic Literature (≤ 2018)

Fasenmyer (1947, 1949); Rainville (1960); Verbaeten (1974); Gosper (1978); Lipshitz (1988); Zeilberger (1982, 1990, 1991); Takayama (1990); Almkvist, Zeilberger (1990); Wilf, Zeilberger (1992); Horneegger (1992); Koornwinder (1993); Paule, Schorn (1995); Majewicz (1996); Riese (1996); Petkovšek, Wilf, Zeilberger (1996); Paule, Riese (1997); Wegschaider (1997); Chyzak, Salvy (1998); Sturmfels, Takayama (1998); Chyzak (2000); Saito, Sturmfels, Takayama (2000); Oaku, Takayama (2001); Le (2001); Riese (2001); Tefera (2000, 2002); Riese (2003); Apagodu, Zeilberger (2006); Kauers (2007); Chen W.Y.C., Sun (2009); Chyzak, Kauers, Salvy (2009); Koutschan (2010); Bostan, Chen S., Chyzak, Li (2010); Chen S., Kauers, Singer (2012); Bostan, Lairez, Salvy (2013); Bostan, Chen S., Chyzak, Li, Xin (2013); Chen S., Huang, Kauers, Li (2015); Lairez (2016); Chen S., Kauers, Koutschan (2016); Huang (2016); Bostan, Dumont, Salvy (2016); Hoeven (2017); Chen S., Hoeij, Kauers, Koutschan (2018); Bostan, Chyzak, Lairez, Salvy (2018).

Applicable to

first-order vs higher-order equations; shift vs differential vs q -analogues vs mixed; ∂ -finite vs non- ∂ -finite; w/ vs wo/ certificate.

Approaches

- bound denominators + bound degrees + linear algebra
- bound denominators + solve functional equations
- elimination by skew Gröbner bases/skew resultants
- reduction of singularity order + linear algebra

History of Algorithms for Creative Telescoping

Algorithmic Literature (≤ 2018)

Fasenmyer (1947, 1949); Rainville (1960); Verbaeten (1974); Gosper (1978); Lipshitz (1988); Zeilberger (1982, 1990, 1991); Takayama (1990); Almkvist, Zeilberger (1990); Wilf, Zeilberger (1992); Hornegger (1992); Koornwinder (1993); Paule, Schorn (1995); Majewicz (1996); Riese (1996); Petkovšek, Wilf, Zeilberger (1996); Paule, Riese (1997); Wegschaider (1997); Chyzak, Salvy (1998); Sturmfels, Takayama (1998); Chyzak (2000); Saito, Sturmfels, Takayama (2000); Oaku, Takayama (2001); Le (2001); Riese (2001); Tefera (2000, 2002); Riese (2003); Apagodu, Zeilberger (2006); Kauers (2007); Chen W.Y.C., Sun (2009); Chyzak, Kauers, Salvy (2009); Koutschan (2010); Bostan, Chen S., Chyzak, Li (2010); Chen S., Kauers, Singer (2012); Bostan, Lairez, Salvy (2013); Bostan, Chen S., Chyzak, Li, Xin (2013); Chen S., Huang, Kauers, Li (2015); Lairez (2016); Chen S., Kauers, Koutschan (2016); Huang (2016); Bostan, Dumont, Salvy (2016); Hoeven (2017); Chen S., Hoeij, Kauers, Koutschan (2018); Bostan, Chyzak, Lairez, Salvy (2018).

Applicable to

first-order vs higher-order equations; shift vs differential vs q -analogues vs mixed; ∂ -finite vs non- ∂ -finite; w/ vs wo/ certificate.

Approaches

- bound denominators + bound degrees + linear algebra
- bound denominators + solve functional equations
- elimination by skew Gröbner bases/skew resultants
- reduction of singularity order + linear algebra

History of Algorithms for Creative Telescoping

Algorithmic Literature (≤ 2018)

Fasenmyer (1947, 1949); Rainville (1960); Verbaeten (1974); Gosper (1978); Lipshitz (1988); Zeilberger (1982, 1990, 1991); Takayama (1990); Almkvist, Zeilberger (1990); Wilf, Zeilberger (1992); Hornegger (1992); Koornwinder (1993); Paule, Schorn (1995); Majewicz (1996); Riese (1996); Petkovšek, Wilf, Zeilberger (1996); Paule, Riese (1997); Wegschaider (1997); Chyzak, Salvy (1998); Sturmfels, Takayama (1998); Chyzak (2000); Saito, Sturmfels, Takayama (2000); Oaku, Takayama (2001); Le (2001); Riese (2001); Tefera (2000, 2002); Riese (2003); Apagodu, Zeilberger (2006); Kauers (2007); Chen W.Y.C., Sun (2009); Chyzak, Kauers, Salvy (2009); Koutschan (2010); Bostan, Chen S., Chyzak, Li (2010); Chen S., Kauers, Singer (2012); Bostan, Lairez, Salvy (2013); Bostan, Chen S., Chyzak, Li, Xin (2013); Chen S., Huang, Kauers, Li (2015); Lairez (2016); Chen S., Kauers, Koutschan (2016); Huang (2016); Bostan, Dumont, Salvy (2016); Hoeven (2017); Chen S., Hoeij, Kauers, Koutschan (2018); Bostan, Chyzak, Lairez, Salvy (2018).

Applicable to

first-order vs higher-order equations; shift vs differential vs q -analogues vs mixed; ∂ -finite vs non- ∂ -finite; w/ vs wo/ certificate.

Approaches

- bound denominators + bound degrees + linear algebra
- bound denominators + solve functional equations
- elimination by skew Gröbner bases/skew resultants
- reduction of singularity order + linear algebra

Running Example

Problem

Integrate $f(n, p, x) = \frac{\exp(-px) T_n(x)}{\sqrt{1-x^2}}$ w.r.t. x and prove the identity

$$F(n, p) := \int_{-1}^{+1} f(n, p, x) dx = (-1)^n \pi I_n(p).$$

Running Example

Problem

Integrate $f(n, p, x) = \frac{\exp(-px) T_n(x)}{\sqrt{1-x^2}}$ w.r.t. x and prove the identity

$$F(n, p) := \int_{-1}^{+1} f(n, p, x) dx = (-1)^n \pi I_n(p).$$

Generating LFEs by algorithm for closure under product yields:

$$\begin{aligned} \frac{\partial f}{\partial p}(n, p, x) + x f(n, p, x) &= 0, \\ n f(n+1, p, x) + (1-x^2) \frac{\partial f}{\partial x}(n, p, x) \\ &\quad + (p(1-x^2) - (n+1)x) f(n, p, x) = 0, \\ (1-x^2) \frac{\partial^2 f}{\partial x^2}(n, p, x) - (2px^2 + 3x - 2p) \frac{\partial f}{\partial x}(n, p, x) \\ &\quad - (p^2 x^2 + 3px - n^2 - p^2 + 1) f(n, p, x) = 0. \end{aligned}$$

Running Example

Problem

Integrate $f(n, p, x) = \frac{\exp(-px) T_n(x)}{\sqrt{1-x^2}}$ w.r.t. x and prove the identity

$$F(n, p) := \int_{-1}^{+1} f(n, p, x) dx = (-1)^n \pi I_n(p).$$

Compact notation using $f_n = f(n+1, p, x)$, $f_x = \frac{\partial f}{\partial x}(n, p, x)$, etc:

$$f_p + xf = 0,$$

$$nf_n + (1-x^2)f_x + (p(1-x^2) - (n+1)x)f = 0,$$

$$(1-x^2)f_{xx} - (2px^2 + 3x - 2p)f_x - (p^2x^2 + 3px - n^2 - p^2 + 1)f = 0.$$

Running Example

Problem

Integrate $f(n, p, x) = \frac{\exp(-px) T_n(x)}{\sqrt{1-x^2}}$ w.r.t. x and prove the identity

$$F(n, p) := \int_{-1}^{+1} f(n, p, x) dx = (-1)^n \pi I_n(p).$$

Compact notation using $f_n = f(n+1, p, x)$, $f_x = \frac{\partial f}{\partial x}(n, p, x)$, etc:

$$f_p + x f = 0,$$

$$n f_n + (1-x^2) f_x + (p(1-x^2) - (n+1)x) f = 0,$$

$$(1-x^2) f_{xx} - (2px^2 + 3x - 2p) f_x - (p^2 x^2 + 3px - n^2 - p^2 + 1) f = 0.$$

Observe: any $f_{n^u p^v x^w}$ is a $\mathbb{Q}(n, p, x)$ -linear combination of f_x and f .

Running Example

Problem

Integrate $f(n, p, x) = \frac{\exp(-px) T_n(x)}{\sqrt{1-x^2}}$ w.r.t. x and prove the identity

$$F(n, p) := \int_{-1}^{+1} f(n, p, x) dx = (-1)^n \pi I_n(p).$$

Compact notation using $f_n = f(n+1, p, x)$, $f_x = \frac{\partial f}{\partial x}(n, p, x)$, etc:

$$f_p + x f = 0,$$

$$n f_n + (1-x^2) f_x + (p(1-x^2) - (n+1)x) f = 0,$$

$$(1-x^2) f_{xx} - (2px^2 + 3x - 2p) f_x - (p^2 x^2 + 3px - n^2 - p^2 + 1) f = 0.$$

Goal: Find a **telescoper** such that there is a **certificate** satisfying

$$\sum_{u,v} c_{u,v}(n, p) f_n^u p^v = g_x \quad \text{for} \quad g = b(n, p, x) f_x + a(n, p, x) f.$$

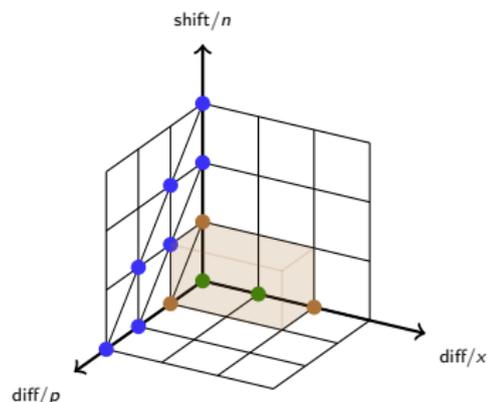
Chyzak's Algorithm (2000): an Example

$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$f_p = (\dots) f_x + (\dots) f,$$

$$f_n = (\dots) f_x + (\dots) f,$$

$$f_{xx} = (\dots) f_x + (\dots) f.$$



For $r = 1, 2, \dots$ until a nonzero equation can be found, solve:

$$\sum_{u+v \leq r} c_{u,v}(n, p) f_n^u p^v = \frac{\partial}{\partial x} (b(n, p, x) f_x + a(n, p, x) f).$$

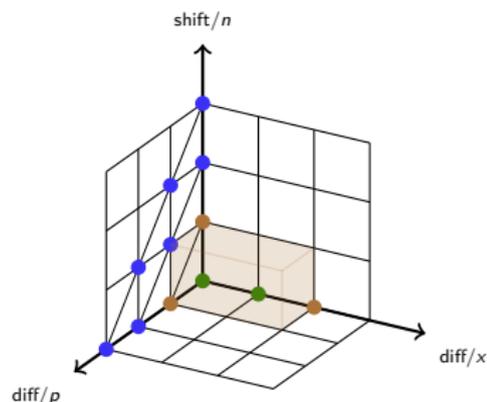
Chyzak's Algorithm (2000): an Example

$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$f_p = (\dots) f_x + (\dots) f,$$

$$f_n = (\dots) f_x + (\dots) f,$$

$$f_{xx} = (\dots) f_x + (\dots) f.$$



For $r = 1, 2, \dots$ until a nonzero equation can be found, solve:

$$\sum_{u+v \leq r} (\dots) c_{u,v}(n, p) f_x^u + (\dots) c_{u,v}(n, p) f = \frac{\partial}{\partial x} (b(n, p, x) f_x + a(n, p, x) f).$$

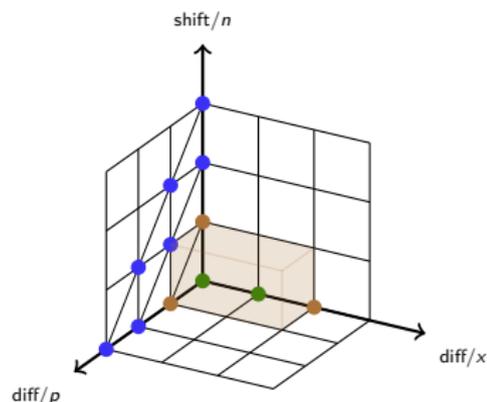
Chyzak's Algorithm (2000): an Example

$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$f_p = (\dots) f_x + (\dots) f,$$

$$f_n = (\dots) f_x + (\dots) f,$$

$$f_{xx} = (\dots) f_x + (\dots) f.$$



For $r = 1, 2, \dots$ until a nonzero equation can be found, solve:

$$\sum_{u+v \leq r} (\dots) c_{u,v} f_x + (\dots) c_{u,v} f = ((\dots) b + b_x + a) f_x + ((\dots) b + a_x) f.$$

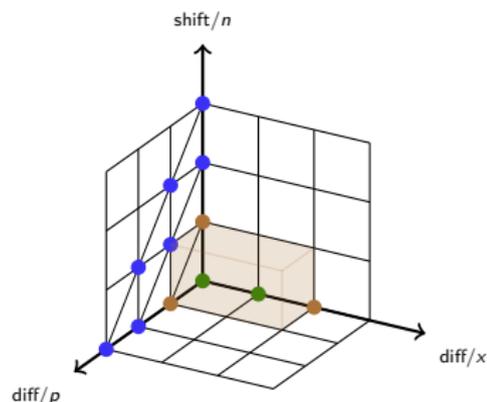
Chyzak's Algorithm (2000): an Example

$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$f_p = (\dots)f_x + (\dots)f,$$

$$f_n = (\dots)f_x + (\dots)f,$$

$$f_{xx} = (\dots)f_x + (\dots)f.$$



For $r = 1, 2, \dots$ until a nonzero equation can be found, solve:

$$\sum_{u+v \leq r} (\dots)c_{u,v} = (\dots)b + b_x + a \quad \text{and} \quad \sum_{u+v \leq r} (\dots)c_{u,v} = (\dots)b + a_x.$$

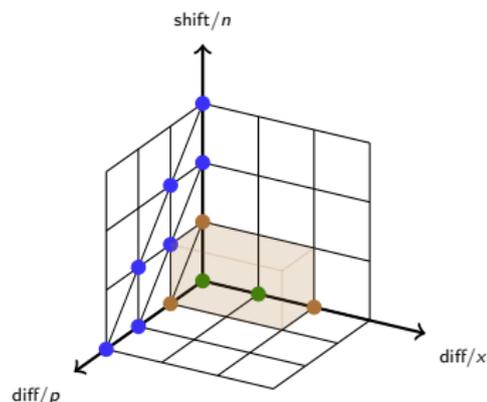
Chyzak's Algorithm (2000): an Example

$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$f_p = (\dots)f_x + (\dots)f,$$

$$f_n = (\dots)f_x + (\dots)f,$$

$$f_{xx} = (\dots)f_x + (\dots)f.$$



For $r = 1, 2, \dots$ until a nonzero equation can be found:

- eliminating a yields: $b_{xx} + (\dots)b_x + (\dots)b = \sum_{u+v \leq r} (\dots)c_{u,v}$;
- a variant of Abramov's decision algorithm finds $b \in \mathbb{Q}(n, p, x)$ and the $c_{u,v} \in \mathbb{Q}(n, p)$; substituting next gives a .

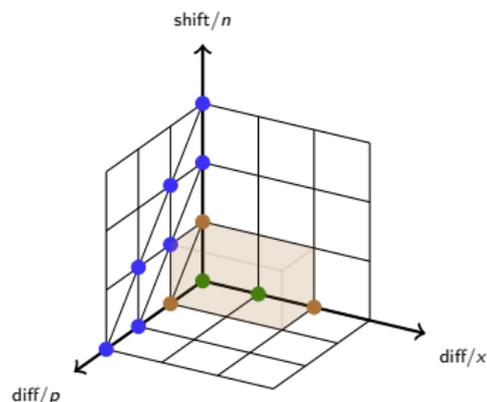
Chyzak's Algorithm (2000): an Example

$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$f_p = (\dots) f_x + (\dots) f,$$

$$f_n = (\dots) f_x + (\dots) f,$$

$$f_{xx} = (\dots) f_x + (\dots) f.$$



For the running example, the algorithm stops at $r = 2$ and outputs:

$$pf_p + pf_n - nf = g_x \quad \text{for} \quad ng = (1 - x^2)f_x + (p(1 - x^2) - x)f,$$

$$pf_{nn} - 2(n + 1)f_n - pf = g_x \quad \text{for}$$

$$ng = 2x(1 - x^2)f_x + 2((px + n)(1 - x^2) - x^2)f.$$

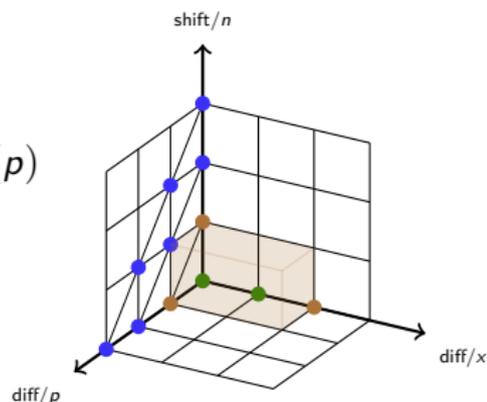
Chyzak's Algorithm (2000): an Example

$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = (-1)^n \pi I_n(p)$$

$$f_p = (\dots) f_x + (\dots) f,$$

$$f_n = (\dots) f_x + (\dots) f,$$

$$f_{xx} = (\dots) f_x + (\dots) f.$$



Upon integrating and using properties of $T_n(\pm 1)$:

$$pF_p + pF_n - nF = [g]_{x=-1}^{x=+1} = 0,$$

$$pF_{nn} - 2(n+1)F_n - pF = [g]_{x=-1}^{x=+1} = 0.$$

One recognizes the equations for the right-hand side $(-1)^n \pi I_n(p)$.

```

[chyzak@slowfox (16:08:44) ~]$ maple -b Mgfun.mla -B
|\~/|      Maple 2018 (X86 64 LINUX)
._\|\|  |/|_ Copyright (c) Maplesoft, a division of Waterloo Maple Inc. 2018
 \ MAPLE / All rights reserved. Maple is a trademark of
 <____ ____> Waterloo Maple Inc.
      |      Type ? for help.
> with(Mgfun);
[MG_Internals, creative_telescoping, dfinite_expr_to_diffeq,

dfinite_expr_to_rec, dfinite_expr_to_sys, diag_of_sys, int_of_sys,

pol_to_sys, rational_creative_telescoping, sum_of_sys, sys*sys, sys+sys]

> f := ChebyshevT(n,x)/sqrt(1-x^2)*exp(-p*x);
      ChebyshevT(n, x) exp(-p x)
      f := -----
              2      1/2
            (-x + 1)

> ct := creative_telescoping(f, {n::shift, p::diff}, x::diff);
memory used=30.3MB, alloc=78.3MB, time=0.37
ct := [[p _F(n + 1, p) + p |-- _F(n, p)| - n _F(n, p),
      \dp                /

x _f(n, p, x) - _f(n + 1, p, x)], [

-p _F(n, p) + p _F(n + 2, p) + (-2 n - 2) _F(n + 1, p),

-2 x _f(n + 1, p, x) + 2 _f(n, p, x)]]

```

Chyzak's Algorithm: Three Issues

- 1 The telescoper (wanted output) is a by-product of the certificate, which is obtained in dense, expanded form (likely to be unneeded in further calculations).
- 2 In dense, expanded form, the certificate is intrinsically large.
- 3 The rational-solving step is sensitive to r , allowing for little reuse of intermediate calculations.

Chyzak's Algorithm: Three Issues

- 1 The telescoper (wanted output) is a by-product of the certificate, which is obtained in dense, expanded form (likely to be unneeded in further calculations).
- 2 In dense, expanded form, the certificate is intrinsically large.
- 3 The rational-solving step is sensitive to r , allowing for little reuse of intermediate calculations.

Example (walks in \mathbb{N}^2 using $\nwarrow, \leftarrow, \downarrow, \rightarrow, \nearrow$, counted by length):

$$\oint \oint \frac{-(1+x)(1+x^2-xy^2)}{(1+x^2)(1-y)(t-xy+ty+tx^2+tx^2y+txy^2)} dx dy$$

$$(16312320t^{20} + \dots)f_{t^5} + (407808000t^{19} + \dots)f_{t^4} + \dots = \frac{\partial g}{\partial x} + \frac{\partial h}{\partial y}$$

LHS = 2 kB, $g = 33$ kB, $h = 896$ kB

Rational Integration: the Classics

Hermite reduction (Ostrogradsky, 1833, 1844/45; Hermite, 1872)

Given P/Q , Hermite reduction finds polynomials A and a such that

$$\int \frac{P(x)}{Q(x)} dx = \frac{A(x)}{Q^-(x)} + \int \frac{a(x)}{Q^*(x)} dx,$$

where $Q^*(x)$ is the squarefree part of $Q(x)$, $Q(x) = Q^-(x)Q^*(x)$, and $\deg a < \deg Q^*$.

Squarefree factorization

Given Q monic, one can in good complexity compute m and 2-by-2 coprime monic Q_i satisfying

$$Q = Q_1^1 Q_2^2 \dots Q_m^m, \quad Q^- = Q_2^1 \dots Q_m^{m-1}, \quad Q^* = Q_1 Q_2 \dots Q_m.$$

Rational Integration: the Classics

Hermite reduction (Ostrogradsky, 1833, 1844/45; Hermite, 1872)

Given P/Q , Hermite reduction finds polynomials A and a such that

$$\int \frac{P(x)}{Q(x)} dx = \frac{A(x)}{Q^-(x)} + \int \frac{a(x)}{Q^*(x)} dx,$$

where $Q^*(x)$ is the squarefree part of $Q(x)$, $Q(x) = Q^-(x)Q^*(x)$, and $\deg a < \deg Q^*$.

Logarithmic part = **obstruction** to existence of rational primitive

For $R(w) = \text{res}_x(b(x), a(x) - b'(x)w)$,

$$\int \frac{a(x)}{b(x)} dx = \sum_{R(c)=0} c \ln(\text{gcd}(b(x), a(x) - b'(x)c))$$

(Trager, 1976).

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Hermite reduction in $K(x)$

Given P/Q , find polynomials A and a with $\deg a < \deg Q^*$ such that

$$\int \frac{P}{Q} dx = \frac{A}{Q^-} + \int \frac{a}{Q^*} dx.$$

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Hermite reduction in $K(x)$

Given P/Q , find polynomials A and a with $\deg a < \deg Q^*$ such that

$$\frac{P}{Q} = \frac{\partial}{\partial x} \left(\frac{A}{Q^-} \right) + \frac{a}{Q^*}.$$

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Bivariate Hermite reduction for creative telescoping in $K(t, x)$

$$\frac{P}{Q} = \frac{\partial}{\partial x} \left(\frac{A^{(0)}}{Q^-} \right) + \frac{a^{(0)}}{Q^*}.$$

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Bivariate Hermite reduction for creative telescoping in $K(t, x)$

$$\frac{P}{Q} = \frac{\partial}{\partial x} \left(\frac{A^{(0)}}{Q^-} \right) + \frac{a^{(0)}}{Q^*}.$$

$$\left(\frac{P}{Q} \right)_t = \frac{\partial}{\partial x} \left(\left(\frac{A^{(0)}}{Q^-} \right)_t \right) + \frac{a_t^{(0)}}{Q^*} - \frac{a^{(0)} Q_t^*}{(Q^*)^2}.$$

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Bivariate Hermite reduction for creative telescoping in $K(t, x)$

$$\frac{P}{Q} = \frac{\partial}{\partial x} \left(\frac{A^{(0)}}{Q^-} \right) + \frac{a^{(0)}}{Q^*}.$$

$$\left(\frac{P}{Q} \right)_t = \frac{\partial}{\partial x} \left(\left(\frac{A^{(0)}}{Q^-} \right)_t \right) + \frac{a_t^{(0)}}{Q^*} + \frac{\partial}{\partial x} \left(\frac{B^{(0)}}{Q^*} \right) + \frac{b^{(0)}}{Q^*}.$$

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Bivariate Hermite reduction for creative telescoping in $K(t, x)$

$$\frac{P}{Q} = \frac{\partial}{\partial x} \left(\frac{A^{(0)}}{Q^-} \right) + \frac{a^{(0)}}{Q^*}.$$

$$\left(\frac{P}{Q} \right)_t = \frac{\partial}{\partial x} \left(\left(\frac{A^{(0)}}{Q^-} \right)_t + \frac{B^{(0)}}{Q^*} \right) + \frac{a_t^{(0)} + b^{(0)}}{Q^*}.$$

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Bivariate Hermite reduction for creative telescoping in $K(t, x)$

$$\frac{P}{Q} = \frac{\partial}{\partial x} \left(\frac{A^{(0)}}{Q^-} \right) + \frac{a^{(0)}}{Q^*}.$$

$$\left(\frac{P}{Q} \right)_t = \frac{\partial}{\partial x} \left(E^{(1)} \right) + \frac{a^{(1)}}{Q^*}.$$

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Bivariate Hermite reduction for creative telescoping in $K(t, x)$

$$\frac{P}{Q} = \frac{\partial}{\partial x} \left(\frac{A^{(0)}}{Q^-} \right) + \frac{a^{(0)}}{Q^*}.$$

$$\left(\frac{P}{Q} \right)_{t^i} = \frac{\partial}{\partial x} \left(E^{(i)} \right) + \frac{a^{(i)}}{Q^*}.$$

Hermite Revisited (Bostan, Chen, Chyzak, Li, 2010)

$$F(t) := \oint \frac{P(t, x)}{Q(t, x)} dx = ? \quad \text{ODE w.r.t. } t?$$

Bivariate Hermite reduction for creative telescoping in $K(t, x)$

$$\frac{P}{Q} = \frac{\partial}{\partial x} \left(\frac{A^{(0)}}{Q^-} \right) + \frac{a^{(0)}}{Q^*}.$$

$$\left(\frac{P}{Q} \right)_{t^i} = \frac{\partial}{\partial x} \left(E^{(i)} \right) + \frac{a^{(i)}}{Q^*}.$$

- **Confinement** $\deg_x a^{(i)} < d := \deg_x Q^* \leq \deg_x Q$:

$$\sum_{i=0}^d c_i(t) a^{(i)}(t, x) = 0 \implies \sum_{i=0}^d c_i F_{t^i} = 0.$$

- **Incremental algorithm** that does not compute $(P/Q)_{t^i}$.
- Degree bounds in $K(t) + \text{eval./interpol.} \implies$ **good complexity**.

Key Idea: Reduce Coordinates, not Functions

D-finite functions can have **complicated singularities**.
Rational functions have only **poles**.

Previous algorithms

$$f \rightarrow [f]$$

$$f_t \rightarrow [f_t]$$

$$f_{tt} \rightarrow [f_{tt}]$$

$$\vdots$$

New algorithm

$$f = R_0 f \rightarrow [R_0] f$$

$$f_t = R_1 f \rightarrow [R_1] f$$

$$f_{tt} = R_2 f \rightarrow [R_2] f$$

$$\vdots$$

Operator Notation

Algebra of linear differential operators with rational coefficients

$$\begin{aligned} \mathbb{A} &= K(t, x) \langle \partial_t, \partial_x \rangle, & M_f &= \mathbb{A}(f) = \{P(f) : P \in \mathbb{A}\} \\ P = \sum p_{i,j}(t, x) \partial_t^i \partial_x^j \in \mathbb{A} &\implies P(f) = \sum p_{i,j}(t, x) f_{t^i x^j} \in M_f \\ \mathbb{S} &= K(t, x) \langle \partial_x \rangle \subset \mathbb{A} \end{aligned}$$

Operator Notation

Algebra of linear differential operators with rational coefficients

$$\begin{aligned} \mathbb{A} &= K(t, x) \langle \partial_t, \partial_x \rangle, & M_f &= \mathbb{A}(f) = \{P(f) : P \in \mathbb{A}\} \\ P &= \sum p_{i,j}(t, x) \partial_t^i \partial_x^j \in \mathbb{A} \implies P(f) = \sum p_{i,j}(t, x) f_{t^i x^j} \in M_f \\ \mathbb{S} &= K(t, x) \langle \partial_x \rangle \subset \mathbb{A} \end{aligned}$$

Hypotheses of D-finiteness

- f is D-finite w.r.t. $\mathbb{A} \implies d := \dim_{K(t,x)}(M_f) < \infty$.
- Let $h \in M_f$ be cyclic, that is to say, $M_f = \bigoplus_{i=0}^{d-1} K(t, x) h_{x^i} = \mathbb{S}(h)$.
- For all $g \in M_f$, there is $A_g \in \mathbb{S}$ such that $g = A_g(h)$.

Operator Notation

Algebra of linear differential operators with rational coefficients

$$\begin{aligned} \mathbb{A} &= K(t, x) \langle \partial_t, \partial_x \rangle, & M_f &= \mathbb{A}(f) = \{P(f) : P \in \mathbb{A}\} \\ P &= \sum p_{i,j}(t, x) \partial_t^i \partial_x^j \in \mathbb{A} \implies P(f) = \sum p_{i,j}(t, x) f_{t^i x^j} \in M_f \\ \mathbb{S} &= K(t, x) \langle \partial_x \rangle \subset \mathbb{A} \end{aligned}$$

Hypotheses of D-finiteness

- f is D-finite w.r.t. $\mathbb{A} \implies d := \dim_{K(t,x)}(M_f) < \infty$.
- Let $h \in M_f$ be cyclic, that is to say, $M_f = \bigoplus_{i=0}^{d-1} K(t, x) h_{x^i} = \mathbb{S}(h)$.
- For all $g \in M_f$, there is $A_g \in \mathbb{S}$ such that $g = A_g(h)$.

Interpretation of creative telescoping

Given f , find a telescoper $T \in K(t) \langle \partial_t \rangle$ and a certificate $g \in M_f$ such that $T(f) = \partial_x(g)$. This really computes $(K(t) \langle \partial_t \rangle)(f) \cap \partial_x(M_f)$.

Lagrange's Identity

Dual of operators

$$P = \sum_{i=0}^r p_i(t, x) \partial_x^i \in \mathbb{S} \xleftrightarrow{*} P^* = \sum_{i=0}^r (-\partial_x)^i p_i(t, x) \in \mathbb{S}$$

Lagrange's Identity

Dual of operators

$$P = \sum_{i=0}^r p_i(t, x) \partial_x^i \in \mathcal{S} \xleftrightarrow{*} P^* = \sum_{i=0}^r (-\partial_x)^i p_i(t, x) \in \mathcal{S}$$

Lagrange's identity

There is a map \mathcal{L}_P , bilinear w.r.t. $(h, \dots, h_{x(r-1)})$ and $(u, \dots, u_{x(r-1)})$, such that

$$\forall u \in K(t, x), \forall h \in M_f, uP(h) - P^*(u)h = \partial_x(\mathcal{L}_P(h, u)).$$

Lagrange's Identity

Dual of operators

$$P = \sum_{i=0}^r p_i(t, x) \partial_x^i \in \mathcal{S} \xleftrightarrow{*} P^* = \sum_{i=0}^r (-\partial_x)^i p_i(t, x) \in \mathcal{S}$$

Lagrange's identity

There is a map \mathcal{L}_P , bilinear w.r.t. $(h, \dots, h_{x^{(r-1)}})$ and $(u, \dots, u_{x^{(r-1)}})$, such that

$$\forall u \in K(t, x), \forall h \in M_f, uP(h) - P^*(u)h = \partial_x(\mathcal{L}_P(h, u)).$$

Proof:
$$\mathcal{L}_P(h, u) = \sum_{i=0}^r \sum_{j=0}^{i-1} (-1)^j (p_i u)_{x^j} h_{x^{i-j-1}}.$$

Consequences of Lagrange's Identity

Lagrange's identity:

$$\forall h \in M_f, \forall u \in K(t, x), uP(h) - P^*(u)h = \partial_x(\mathcal{L}_P(h, u)).$$

Let h be cyclic and $L \in S$ be such that $L(h) = 0$. Then, for all $g \in M_f$:

Operator to rational function: $\mathbb{A}(f) = S(h) \rightarrow K(t, x)h$

$$g \in A_g^*(1)h + \partial_x(M_f).$$

[by $u = 1, P = A_g$]

Consequences of Lagrange's Identity

Lagrange's identity:

$$\forall h \in M_f, \forall u \in K(t, x), uP(h) - P^*(u)h = \partial_x(\mathcal{L}_P(h, u)).$$

Let h be cyclic and $L \in \mathcal{S}$ be such that $L(h) = 0$. Then, for all $g \in M_f$:

Operator to rational function: $\mathbb{A}(f) = \mathcal{S}(h) \rightarrow K(t, x)h$

$$g \in A_g^*(1)h + \partial_x(M_f). \quad [\text{by } u = 1, P = A_g]$$

Equivalent rational factors: $K(t, x)h \rightarrow (K(t, x) \bmod L^*(K(t, x)))h$

$$\forall u \in K(t, x), g \in (A_g^*(1) - L^*(u))h + \partial_x(M_f). \quad [\text{by } P = L]$$

Consequences of Lagrange's Identity

Lagrange's identity:

$$\forall h \in M_f, \forall u \in K(t, x), uP(h) - P^*(u)h = \partial_x(\mathcal{L}_P(h, u)).$$

Let h be cyclic and $L \in \mathcal{S}$ be such that $L(h) = 0$. Then, for all $g \in M_f$:

Operator to rational function: $\mathbb{A}(f) = \mathcal{S}(h) \rightarrow K(t, x)h$

$$g \in A_g^*(1)h + \partial_x(M_f). \quad [\text{by } u = 1, P = A_g]$$

Equivalent rational factors: $K(t, x)h \rightarrow (K(t, x) \bmod L^*(K(t, x)))h$

$$\forall u \in K(t, x), g \in (A_g^*(1) - L^*(u))h + \partial_x(M_f). \quad [\text{by } P = L]$$

Testing derivatives (for L of minimal order)

$$g \in \partial_x(M_f) \Rightarrow \exists q \in K(t, x), A_g^*(1) = L^*(q). \quad [\text{by } A_g^*(1) \in \partial_x \mathcal{S} + \mathcal{S}L]$$

Consequences of Lagrange's Identity

Lagrange's identity:

$$\forall h \in M_f, \forall u \in K(t, x), uP(h) - P^*(u)h = \partial_x(\mathcal{L}_P(h, u)).$$

Let h be cyclic and $L \in \mathbb{S}$ be such that $L(h) = 0$. Then, for all $g \in M_f$:

Operator to rational function: $\mathbb{A}(f) = \mathbb{S}(h) \rightarrow K(t, x)h$

$$g \in A_g^*(1)h + \partial_x(M_f).$$

Equivalent rational factors: $K(t, x)h \rightarrow (K(t, x) \bmod L^*(K(t, x)))h$

$$\forall u \in K(t, x), g \in (A_g^*(1) - L^*(u))h + \partial_x(M_f).$$

Testing derivatives (for L of minimal order)

$$g \in \partial_x(M_f) \Leftrightarrow \exists q \in K(t, x), A_g^*(1) = L^*(q).$$

Consequences of Lagrange's Identity

Lagrange's identity:

$$\forall h \in M_f, \forall u \in K(t, x), uP(h) - P^*(u)h = \partial_x(\mathcal{L}_P(h, u)).$$

Let h be cyclic and $L \in \mathbb{S}$ be such that $L(h) = 0$. Then, for all $g \in M_f$:

Operator to rational function: $\mathbb{A}(f) = \mathbb{S}(h) \rightarrow K(t, x)h$

$$g \in A_g^*(1)h + \partial_x(M_f).$$

Equivalent rational factors: $K(t, x)h \rightarrow (K(t, x) \bmod L^*(K(t, x)))h$

$$\forall u \in K(t, x), g \in (A_g^*(1) - L^*(u))h + \partial_x(M_f). \quad [\text{Reduction?}]$$

Testing derivatives (for L of minimal order)

$$g \in \partial_x(M_f) \Leftrightarrow \exists q \in K(t, x), A_g^*(1) = L^*(q). \quad [\text{Algorithm?}]$$

Running Example (continued)

Problem

Integrate $f(n, p, x) = \frac{\exp(-px) T_n(x)}{\sqrt{1-x^2}}$ w.r.t. x and prove the identity

$$F(n, p) := \int_{-1}^{+1} f(n, p, x) dx = (-1)^n \pi I_n(p).$$

In operator notation, f is cancelled by **all left-linear combinations** of:

$$\begin{aligned} & \partial_p + x1, \quad n\partial_n + (1-x^2)\partial_x + (p(1-x^2) - (n+1)x)1, \\ & (1-x^2)\partial_x^2 - (2px^2 + 3x - 2p)\partial_x - (p^2x^2 + 3px - n^2 - p^2 + 1)1. \end{aligned}$$

Goal: Find a **telescoper** such that there is a **certificate** satisfying

$$\sum_{u,v} c_{u,v}(n, p) \partial_n^u \partial_p^v = \partial_x (b(n, p, x) \partial_x + a(n, p, x) 1)$$

modulo the operators above.

Running Example (continued)

Problem

Integrate $f(n, p, x) = \frac{\exp(-px) T_n(x)}{\sqrt{1-x^2}}$ w.r.t. x and prove the identity

$$F(n, p) := \int_{-1}^{+1} f(n, p, x) dx = (-1)^n \pi I_n(p).$$

In operator notation, f is cyclic, so $h := f$, and it is cancelled by:

$$L := \partial_p + x1, \quad n\partial_n + (1-x^2)\partial_x + (p(1-x^2) - (n+1)x)1,$$

$$L := (1-x^2)\partial_x^2 - (2px^2 + 3x - 2p)\partial_x - (p^2x^2 + 3px - n^2 - p^2 + 1)1.$$

Goal: Find a **telescoper** such that there is a **certificate** satisfying

$$\sum_{u,v} c_{u,v}(n, p) \partial_n^u \partial_p^v = \partial_x (b(n, p, x) \partial_x + a(n, p, x) 1)$$

modulo the operators above.

Reduction-Based CT Algorithm (2018): an Example

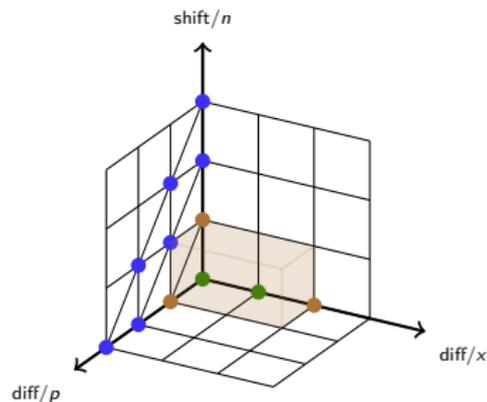
$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$L_1 := \partial_p - (\dots)\partial_x - (\dots)1,$$

$$L_2 := \partial_n - (\dots)\partial_x - (\dots)1,$$

$$L_3 := \partial_x^2 - (\dots)\partial_x - (\dots)1.$$

$$L := L_3, \quad I := \mathbb{A}L_1 + \mathbb{A}L_2 + \mathbb{A}L_3.$$



Reduction-Based CT Algorithm (2018): an Example

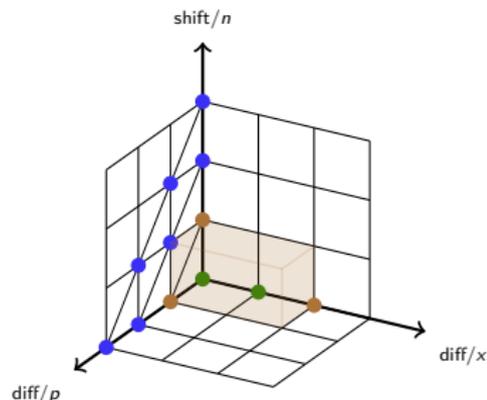
$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$L_1 := \partial_p - (\dots)\partial_x - (\dots)\mathbf{1},$$

$$L_2 := \partial_n - (\dots)\partial_x - (\dots)\mathbf{1},$$

$$L_3 := \partial_x^2 - (\dots)\partial_x - (\dots)\mathbf{1}.$$

$$L := L_3, \quad I := \mathbb{A}L_1 + \mathbb{A}L_2 + \mathbb{A}L_3.$$



For $P = 1, \partial_n, \partial_p, \partial_n^2, \partial_n\partial_p, \partial_p^2$:

- set $g = P(h)$, so that $A_g = \text{rem}(P, I) = v(p, n, x)\partial_x + u(p, n, x)\mathbf{1}$,

Reduction-Based CT Algorithm (2018): an Example

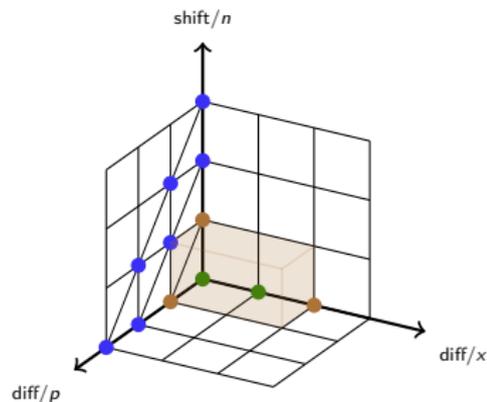
$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$L_1 := \partial_p - (\dots)\partial_x - (\dots)\mathbf{1},$$

$$L_2 := \partial_n - (\dots)\partial_x - (\dots)\mathbf{1},$$

$$L_3 := \partial_x^2 - (\dots)\partial_x - (\dots)\mathbf{1}.$$

$$L := L_3, \quad I := \mathbb{A}L_1 + \mathbb{A}L_2 + \mathbb{A}L_3.$$



For $P = 1, \partial_n, \partial_p, \partial_n^2, \partial_n\partial_p, \partial_p^2$:

- set $g = P(h)$, so that $A_g = \text{rem}(P, I) = v(p, n, x)\partial_x + u(p, n, x)\mathbf{1}$,
- $A_g^* = -v\partial_x + (u - v_x)$, so that $g = (u - v_x)f + \partial_x(\dots)$.

Reduction-Based CT Algorithm (2018): an Example

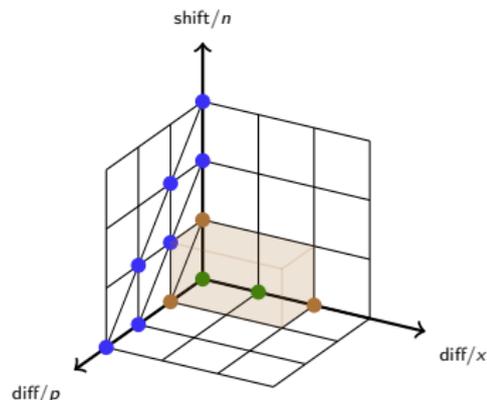
$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$L_1 := \partial_p - (\dots)\partial_x - (\dots)1,$$

$$L_2 := \partial_n - (\dots)\partial_x - (\dots)1,$$

$$L_3 := \partial_x^2 - (\dots)\partial_x - (\dots)1.$$

$$L := L_3, \quad I := \mathbb{A}L_1 + \mathbb{A}L_2 + \mathbb{A}L_3.$$



For $P = 1, \partial_n, \partial_p, \partial_n^2, \partial_n\partial_p, \partial_p^2$:

- set $g = P(h)$, so that $A_g = \text{rem}(P, I) = v(p, n, x)\partial_x + u(p, n, x)1$,
- $A_g^* = -v\partial_x + (u - v_x)$, so that $g = (u - v_x)f + \partial_x(\dots)$.

For those $P, u - v_x \in K(p, n)[x]$ with degree ≤ 3 , while

$$L^*(p^2x^0) = p^2x^2 - px - (n^2 + p^2),$$

$$L^*(p^2x^1) = p^2x^3 - 3px^2 - (n^2 + p^2 - 1)x + 2p.$$

Reduction-Based CT Algorithm (2018): an Example

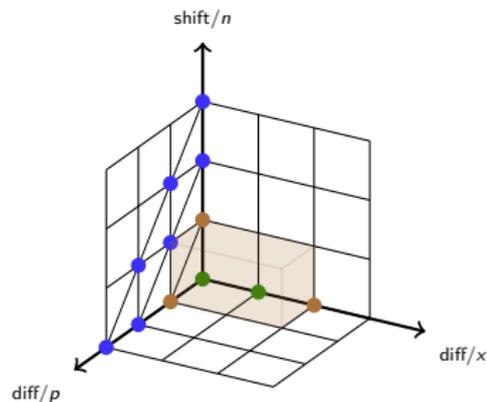
$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$L_1 := \partial_p - (\dots)\partial_x - (\dots)1,$$

$$L_2 := \partial_n - (\dots)\partial_x - (\dots)1,$$

$$L_3 := \partial_x^2 - (\dots)\partial_x - (\dots)1.$$

$$L := L_3, \quad I := \mathbb{A}L_1 + \mathbb{A}L_2 + \mathbb{A}L_3.$$



For $P = 1, \partial_n, \partial_p, \partial_n^2, \partial_n\partial_p, \partial_p^2$:

$$P(f) = (u - v_x)f + \partial_x(\dots) = (\mu_P(p, n)x^1 + \lambda_P(p, n)x^0) f + \partial_x(\dots).$$

Reduction-Based CT Algorithm (2018): an Example

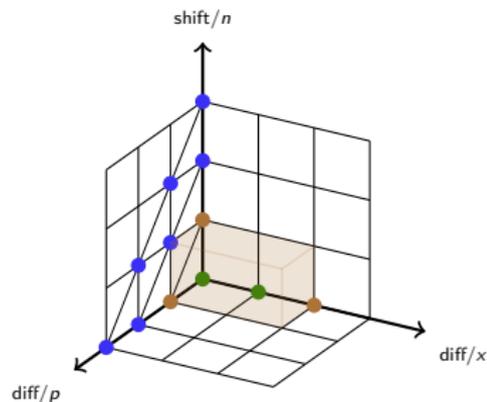
$$\int_{-1}^{+1} f(n, p, x) dx = F(n, p) = ?$$

$$L_1 := \partial_p - (\dots)\partial_x - (\dots)1,$$

$$L_2 := \partial_n - (\dots)\partial_x - (\dots)1,$$

$$L_3 := \partial_x^2 - (\dots)\partial_x - (\dots)1.$$

$$L := L_3, \quad I := \mathbb{A}L_1 + \mathbb{A}L_2 + \mathbb{A}L_3.$$



For $P = 1, \partial_n, \partial_p, \partial_n^2, \partial_n\partial_p, \partial_p^2$:

$$P(f) = (u - v_x)f + \partial_x(\dots) = (\mu_P(p, n)x^1 + \lambda_P(p, n)x^0) f + \partial_x(\dots).$$

Linear algebra over $K(p, n)$ finds a basis of **telescopers**

$$\left(\sum_P c_P P \right) (f) = \partial_x(\dots).$$

Reduction Modulo $L^*(K(x))$

Local decomposition of a rational function $R \in K(x)$

$$R = R_{(\infty)} + \sum_{\alpha} R_{(\alpha)} \text{ for some } R_{(\alpha)} \in \frac{1}{x-\alpha} K(\alpha) \left[\frac{1}{x-\alpha} \right] \text{ and } R_{(\infty)} \in K[x].$$

Reduction Modulo $L^*(K(x))$

Local decomposition of a rational function $R \in K(x)$

$$R = R_{(\infty)} + \sum_{\alpha} R_{(\alpha)} \text{ for some } R_{(\alpha)} \in \frac{1}{x-\alpha} K(\alpha) \left[\frac{1}{x-\alpha} \right] \text{ and } R_{(\infty)} \in K[x].$$

Local study of the action of L^*

\exists polynomials I_{α} and I_{∞} , \exists integers σ_{α} and σ_{∞} , such that $\forall s \in \mathbb{Z}$,

$$L^*((x-\alpha)^{-s}) \underset{x \rightarrow \alpha}{=} I_{\alpha}(-s)(x-\alpha)^{\sigma_{\alpha}-s} + \mathcal{O}((x-\alpha)^{\sigma_{\alpha}-(s-1)}),$$

$$L^*((1/x)^{-s}) \underset{x \rightarrow \infty}{=} I_{\infty}(-s)(1/x)^{\sigma_{\infty}-s} + \mathcal{O}((1/x)^{\sigma_{\infty}-(s-1)}).$$

Reduction Modulo $L^*(K(x))$

Local decomposition of a rational function $R \in K(x)$

$$R = R_{(\infty)} + \sum_{\alpha} R_{(\alpha)} \text{ for some } R_{(\alpha)} \in \frac{1}{x-\alpha} K(\alpha) \left[\frac{1}{x-\alpha} \right] \text{ and } R_{(\infty)} \in K[x].$$

Local study of the action of L^*

\exists polynomials l_{α} and l_{∞} , \exists integers σ_{α} and σ_{∞} , such that $\forall s \in \mathbb{Z}$,

$$L^* \left((x-\alpha)^{-s-\sigma_{\alpha}} \right) \underset{x \rightarrow \alpha}{=} l_{\alpha}(-s-\sigma_{\alpha})(x-\alpha)^{-s} + \mathcal{O}((x-\alpha)^{-(s-1)}),$$

$$L^* \left((1/x)^{-s-\sigma_{\infty}} \right) \underset{x \rightarrow \infty}{=} l_{\infty}(-s-\sigma_{\infty})(1/x)^{-s} + \mathcal{O}((1/x)^{-(s-1)}).$$

Reduction Modulo $L^*(K(x))$

Local decomposition of a rational function $R \in K(x)$

$$R = R_{(\infty)} + \sum_{\alpha} R_{(\alpha)} \text{ for some } R_{(\alpha)} \in \frac{1}{x-\alpha} K(\alpha) \left[\frac{1}{x-\alpha} \right] \text{ and } R_{(\infty)} \in K[x].$$

Local study of the action of L^*

\exists polynomials l_{α} and l_{∞} , \exists integers σ_{α} and σ_{∞} , such that $\forall s \in \mathbb{Z}$,

$$(x - \alpha)^{-s} \underset{x \rightarrow \alpha}{=} l_{\alpha} (-s - \sigma_{\alpha})^{-1} L^* \left((x - \alpha)^{-s - \sigma_{\alpha}} \right) + \mathcal{O} \left((x - \alpha)^{-(s-1)} \right),$$

$$(1/x)^{-s} \underset{x \rightarrow \infty}{=} l_{\infty} (-s - \sigma_{\infty})^{-1} L^* \left((1/x)^{-s - \sigma_{\infty}} \right) + \mathcal{O} \left((1/x)^{-(s-1)} \right).$$

Reduction Modulo $L^*(K(x))$

Local decomposition of a rational function $R \in K(x)$

$$R = R_{(\infty)} + \sum_{\alpha} R_{(\alpha)} \text{ for some } R_{(\alpha)} \in \frac{1}{x-\alpha} K(\alpha) \left[\frac{1}{x-\alpha} \right] \text{ and } R_{(\infty)} \in K[x].$$

Local study of the action of L^*

\exists polynomials l_{α} and l_{∞} , \exists integers σ_{α} and σ_{∞} , such that $\forall s \in \mathbb{Z}$,

$$(x - \alpha)^{-s} \underset{x \rightarrow \alpha}{=} l_{\alpha}(-s - \sigma_{\alpha})^{-1} L^* \left((x - \alpha)^{-s - \sigma_{\alpha}} \right) + \mathcal{O} \left((x - \alpha)^{-(s-1)} \right),$$

$$(1/x)^{-s} \underset{x \rightarrow \infty}{=} l_{\infty}(-s - \sigma_{\infty})^{-1} L^* \left((1/x)^{-s - \sigma_{\infty}} \right) + \mathcal{O} \left((1/x)^{-(s-1)} \right).$$

Weak reduction strategy

- reduce at finite α (in any order) before at ∞ ,
- skip monomials for which $l_{\alpha}(-s - \sigma_{\alpha}) = 0$ or $l_{\infty}(-s - \sigma_{\infty}) = 0$.

Canonical Form Modulo $L^*(K(x))$

Problem: $L^*(K(x))$ does not weakly reduce to $\{0\}$

For $c_0 = I_\alpha(-s - \sigma_\alpha)$ and some c_1 , write $R := L^*((x - \alpha)^{-s - \sigma_\alpha})$ as

$$R = c_0(x - \alpha)^{-s} + c_1(x - \alpha)^{-(s-1)} + \mathcal{O}((x - \alpha)^{-(s-2)}).$$

- If $c_0 \neq 0$, this reduces to

$$L^*((x - \alpha)^{-s - \sigma_\alpha} - (x - \alpha)^{-s - \sigma_\alpha}) = 0.$$

- If $c_0 = 0$ and $c_1 \neq 0$, this reduces to some

$$L^*\left((x - \alpha)^{-s - \sigma_\alpha} - \frac{c_1}{c_2}(x - \alpha)^{-(s-1) - \sigma_\alpha}\right),$$

which is unlikely to further reduce to 0.

Canonical Form Modulo $L^*(K(x))$

Problem: $L^*(K(x))$ does not weakly reduce to $\{0\}$

For $c_0 = l_\alpha(-s - \sigma_\alpha)$ and some c_1 , write $R := L^*((x - \alpha)^{-s - \sigma_\alpha})$ as

$$R = c_0(x - \alpha)^{-s} + c_1(x - \alpha)^{-(s-1)} + \mathcal{O}((x - \alpha)^{-(s-2)}).$$

- If $c_0 \neq 0$, this reduces to

$$L^*((x - \alpha)^{-s - \sigma_\alpha} - (x - \alpha)^{-s - \sigma_\alpha}) = 0.$$

- If $c_0 = 0$ and $c_1 \neq 0$, this reduces to some

$$L^*\left((x - \alpha)^{-s - \sigma_\alpha} - \frac{c_1}{c_2}(x - \alpha)^{-(s-1) - \sigma_\alpha}\right),$$

which is unlikely to further reduce to 0.

Solution

- finitely-many potential obstructions, described by the integer zeros of the l_α and l_∞ ,
- this can be computed, leading to a canonical-form computation.

```
[chyzak@slowfox (04:21:54) ~]$ maple -b Mgfuns.mla -B
|\~/|      Maple 2018 (X86 64 LINUX)
._|\|  |/|_. Copyright (c) Maplesoft, a division of Waterloo Maple Inc. 2018
\ MAPLE / All rights reserved. Maple is a trademark of
<_---_---> Waterloo Maple Inc.
      |      Type ? for help.
> read "redct.mpl";
> f := ChebyshevT(n,x)/sqrt(1-x^2)*exp(-p*x);
                ChebyshevT(n, x) exp(-p x)
      f := -----
                2      1/2
              (-x  + 1)

> redct(Int(f,x=-1..1),[n::shift,p::diff]);
memory used=3.5MB, alloc=8.3MB, time=0.09
                2
      [p D[n] + p D[p] - n, p D[n]  - 2 n D[n] - p - 2 D[n]]

> f := 2*BesselJ(m+n,2*t*x)*ChebyshevT(m-n,x)/sqrt(1-x^2);
                2 BesselJ(m + n, 2 t x) ChebyshevT(m - n, x)
      f := -----
                2      1/2
              (-x  + 1)

> redct(Int(f,x),[t::diff, n::shift, m::shift]);
memory used=1189.8MB, alloc=144.8MB, time=9.98
                2
      [t D[m] + t D[n] + t D[t] - m - n, t D[m]  - 2 m D[m] + t - 2 D[m],

                2
      t D[n]  - 2 n D[n] + t - 2 D[n]]
```

Timings: More than 140 integrals tested

Algorithm	(1)	(2)	(3)	(4)	(5)	(6)	(7)
new (mpl)	13s	> 1h	> 1h	1.5s	1.5s	165s	53s
Chyzak's (mma)	19s	253s	45s	232s	516s	>1h	>1h
Koutschan's (mma)	1.9s†	2.3s	5.3s	>1h	2.3s†	5.4s	2.2s†

$$\int \frac{{}_2J_{m+n}(2tx)T_{m-n}(x)}{\sqrt{1-x^2}} dx \quad [\text{diff. } t, \text{ shift } n \text{ and } m], \quad (1)$$

$$\int_0^1 C_n^{(\lambda)}(x) C_m^{(\lambda)}(x) C_\ell^{(\lambda)}(x) (1-x^2)^{\lambda-\frac{1}{2}} dx \quad [\text{shift } n, m, \ell], \quad (2)$$

$$\int_0^\infty x J_1(ax) Y_0(x) K_0(x) dx \quad [\text{diff. } a], \quad (3)$$

$$\int \frac{n^2+x+1}{n^2+1} \left(\frac{(x+1)^2}{(x-4)(x-3)^2(x^2-5)^3} \right)^n \sqrt{x^2-5} e^{\frac{x^3+1}{x(x-3)(x-4)^2}} dx \quad (4)$$

[shift n],

$$\int C_m^{(\mu)}(x) C_n^{(\nu)}(x) (1-x^2)^{\nu-1/2} dx \quad [\text{shift } n, m, \mu, \nu], \quad (5)$$

$$\int x^\ell C_m^{(\mu)}(x) C_n^{(\nu)}(x) (1-x^2)^{\nu-1/2} dx \quad [\text{shift } \ell, m, n, \mu, \nu], \quad (6)$$

$$\int (x+a)^{\gamma+\lambda-1} (a-x)^{\beta-1} C_m^{(\gamma)}(x/a) C_n^{(\lambda)}(x/a) dx \quad (7)$$

[diff. a , shift $n, m, \beta, \gamma, \lambda$].

†: Heuristic got these faster answers by looking for telescopers of non-minimal orders, yet smaller sizes.

Timings: More than 140 integrals tested

Algorithm	(1)	(2)	(3)	(4)	(5)	(6)	(7)
new (mpl)	13s	> 1h	> 1h	1.5s	1.5s	165s	53s
Chyzak's (mma)	19s	253s	45s	232s	516s	>1h	>1h
Koutschan's (mma)	1.9s†	2.3s	5.3s	>1h	2.3s†	5.4s	2.2s†

$$\int \frac{2J_{m+n}(2tx)T_{m-n}(x)}{\sqrt{1-x^2}} dx \quad [\text{diff. } t, \text{ shift } n \text{ and } m], \quad (1)$$

$$\int_0^1 C_n^{(\lambda)}(x)C_m^{(\lambda)}(x)C_\ell^{(\lambda)}(x)(1-x^2)^{\lambda-\frac{1}{2}} dx \quad [\text{shift } n, m, \ell], \quad (2)$$

$$\int_0^\infty xJ_1(ax)Y_0(x)K_0(x) dx \quad [\text{diff. } a], \quad (3)$$

$$\int \frac{n^2+x+1}{n^2+1} \left(\frac{(x+1)^2}{(x-4)(x-3)^2(x^2-5)^3} \right)^n \sqrt{x^2-5} e^{\frac{x^3+1}{x(x-3)(x-4)^2}} dx \quad (4)$$

[shift n],

$$\int C_m^{(\mu)}(x)C_n^{(\nu)}(x)(1-x^2)^{\nu-1/2} dx \quad [\text{shift } n, m, \mu, \nu], \quad (5)$$

$$\int x^\ell C_m^{(\mu)}(x)C_n^{(\nu)}(x)(1-x^2)^{\nu-1/2} dx \quad [\text{shift } \ell, m, n, \mu, \nu], \quad (6)$$

$$\int (x+a)^{\gamma+\lambda-1}(a-x)^{\beta-1} C_m^{(\gamma)}(x/a)C_n^{(\lambda)}(x/a) dx \quad (7)$$

[diff. a , shift $n, m, \beta, \gamma, \lambda$].

†: Heuristic got these faster answers by looking for telescopers of non-minimal orders, yet smaller sizes.

Need to investigate failures:

- non-mathematical bugs? “not ours”?
- impact of apparent singularities of P^* ?

Summary

Approach by solving functional equations (1991+)

see failures to solve as obstructions,
recombine obstructions

Primal reduction-based approach (2010+)

work on rational coordinates
to simplify singularities,
Lagrange's formula

Dual reduction-based approach (2018)