

On $(3, 1)$ -Regular Graphs with One More Vertex than Edges: A Case Study in Difference-Differential Algebra

Frédéric Chyzak



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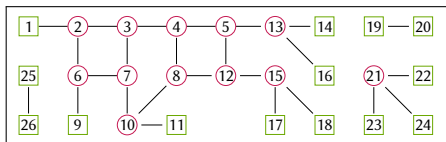
Île d'Oléron, May 24–30, 2026

Joint work with Hui Huang and Manuel Kauers

A Conjecture in Enumerative Graph Theory

Conjecture 16 in (Kauers, Koutschan, 2023)

Sequence A339987 of the OEIS, which counts vertex-labelled $(3,1)$ -regular graphs having one more vertex than edges by half the number of vertices, satisfies an explicit, guessed recurrence relation of order 5, valid for all $k \geq 0$.



26 vertices, 25 edges, all degrees in $\{1, 3\}$

$$a_0 = 0, \quad a_1 = 1, \quad a_2 = 4, \quad a_3 = 90, \quad a_4 = 8400, \quad a_5 = 1426950, \\ a_6 = 366153480, \quad a_7 = 134292027870$$

A Conjecture in Enumerative Graph Theory

$$\begin{aligned} & 32(328k^3 + 3300k^2 + 10844k + 11589) \\ & \quad \times (k+1)(k+2)(2k+1)(2k+3)(2k+5)(2k+7)(2k+9) a_k \\ & - 8(2624k^4 + 30664k^3 + 129460k^2 + 232328k + 148119) \\ & \quad \times (k+2)(2k+3)(2k+5)(2k+7)(2k+9) a_{k+1} \\ & - 16(2952k^5 + 40852k^4 + 219308k^3 + 569267k^2 + 712135k + 341634) \\ & \quad \times (k+3)(2k+5)(2k+7)(2k+9) a_{k+2} \\ & + 8(3936k^5 + 55672k^4 + 306380k^3 + 818282k^2 + 1057879k + 527520) \\ & \quad \times (k+4)(2k+7)(2k+9) a_{k+3} \\ & - 2(2624k^5 + 42472k^4 + 264028k^3 + 786236k^2 + 1117119k + 601452) \\ & \quad \times (k+5)(2k+9) a_{k+4} \\ & + 3(328k^3 + 2316k^2 + 5228k + 3717)(k+4)(k+6) a_{k+5} = 0 \end{aligned}$$

Gussed by computing a Hermite–Padé approximant

generation of enough terms \rightarrow gussed ODE \rightarrow gussed recurrence

Three Proofs by Difference/Differential Algebra

Combinatorial theory: (Gessel, 1987, 1990)

By a representation as a triple-sum

Classical creative telescoping for sums has to be adapted.

Difference (shift) operators with rational coefficients.

(Zeilberger, 1990), (Chyzak, 2000), (Koutschan, 2010)

By a representation as a residue

Uses reduction-based creative telescoping for D-modules.

Differential operators with polynomial coefficients.

(Brochet, Chyzak, Lairez, 2026)

Combinatorial theory: (Read, 1970), (Wormald, 1979)

By combinatorial recursion

Uses differential elimination and creative telescoping for diagonals.

Differential operators with rational coefficients.

(Bostan, Chyzak, Lairez, Salvy, 2018), (van der Hoeven, 2021)

Creative Telescoping for Sums and Integrals

$$U_n = \sum_{k=a}^b u_{n,k} = ?$$

Given a relation $a_r(n)u_{n+r,k} + \cdots + a_0(n)u_{n,k} = v_{n,k+1} - v_{n,k}$, summation leads by “telescoping” to

$$a_r(n)U_{n+r} + \cdots + a_0(n)U_n = v_{n,b+1} - v_{n,a} \stackrel{\text{often}}{=} 0.$$

$$U(t) = \int_a^b u(t, x) dx = ?$$

Given a relation $a_r(t) \frac{\partial^r u}{\partial t^r} + \cdots + a_0(t)u = \frac{\partial}{\partial x} v(t, x)$, integrating leads by “telescoping” to

$$a_r(t) \frac{\partial^r U}{\partial t^r} + \cdots + a_0(t)U = v(t, b) - v(t, a) \stackrel{\text{often}}{=} 0.$$

Adapts easily to $U(t) = \sum_{k=a}^b u_k(t)$, $U_n = \int_a^b u_n(x) dx$, etc.

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Telescoper

“Certificate”

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Older algorithms compute **telescopers** while computing “**certificates**”.
Reduction-based algorithms compute **telescopers** without computing “**certificates**”.

(Chyzak, Salvy, 1998), building on (Ore, 1933)

Difference operators

$\mathbb{Q}(k, u, v, w)\langle S_k, S_u, S_v, S_w \rangle$: skew polynomials in S_k, S_u, S_v, S_w obeying

$$S_k k = (k + 1)S_k, \quad \text{and similar relations,}$$

$$S_k \cdot f_{k,u,v,w} = f_{k+1,u,v,w}, \quad \text{and similar actions.}$$

Also $\mathbb{Q}(k)\langle S_k \rangle$, etc.

Differential operators

$\mathbb{Q}(t)[p_1, p_2, p_3, q]\langle \partial_{p_1}, \partial_{p_2}, \partial_{p_3}, \partial_q \rangle$: skew polynomials in $\partial_{p_1}, \partial_{p_2}, \partial_{p_3}, \partial_q$ obeying

$$\partial_{p_1} p_1 = p_1 \partial_{p_1} + 1, \quad \text{and similar relations,}$$

$$\partial_{p_1} \cdot f(p_1, p_2, p_3, q) = \frac{\partial f}{\partial p_1}(p_1, p_2, p_3, q), \quad \text{and similar actions.}$$

Also $\mathbb{Q}(q, t)\langle \partial_q, \partial_t \rangle$, etc.

First Approach:
By a Triple-Sum Representation

A Formula from the Theory of Symmetric Functions

An application of Gessel's theory (1987, 1990)

$r_{m,n} := \#\{(3, 1)\text{-regular simple graphs on } n \text{ labelled vertices with } m \text{ edges}\}$

The EGF $\sum_{m,n=0}^{\infty} r_{m,n} q^m \frac{t^n}{n!} = \langle F(p, q), G(p, t) \rangle \in \mathbb{Q}[[q, t]],$

where

$$F(p, q) = \exp\left(\frac{1}{2}q(p_1^2 - p_2) - \frac{1}{4}q^2 p_2^2 + \frac{1}{6}q^3 p_3^2\right) \in \mathbb{Q}[p_1, p_2, p_3][[q]],$$

$$G(p, t) = \exp\left((p_1 + \frac{1}{6}p_1^3 + \frac{1}{2}p_1 p_2 + \frac{1}{3}p_3)t\right) \in \mathbb{Q}[p_1, p_2, p_3][[t]],$$

is D-finite w.r.t. q and t .

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Scalar product (= standard pairing on symmetric functions)

A scalar product on $\mathbb{Q}[p_1, p_2, p_3]$ is defined from the formula

$$\langle p_1^{r_1} p_2^{r_2} p_3^{r_3}, p_1^{s_1} p_2^{s_2} p_3^{s_3} \rangle = \begin{cases} 1^{r_1} r_1! 2^{r_2} r_2! 3^{r_3} r_3! & \text{if } r = s, \\ 0 & \text{otherwise.} \end{cases}$$

This extends by bilinearity to $\mathbb{Q}[p_1, p_2, p_3][[q, t]].$

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The EGF $\sum_{m,n=0}^{\infty} r_{m,n} q^m \frac{t^n}{n!} = \langle F(p, q), G(p, t) \rangle \in \mathbb{Q}[[q, t]],$

where

$$F(p, q) = \exp\left(\frac{1}{2}q(p_1^2 - p_2) - \frac{1}{4}q^2 p_2^2 + \frac{1}{6}q^3 p_3^2\right) \in \mathbb{Q}[p_1, p_2, p_3][[q]],$$

$$G(p, t) = \exp\left((p_1 + \frac{1}{6}p_1^3 + \frac{1}{2}p_1 p_2 + \frac{1}{3}p_3)t\right) \in \mathbb{Q}[p_1, p_2, p_3][[t]],$$

is D-finite w.r.t. q and t .

Earlier algorithms for computing scalar products without q

(Chyzak, Mishna, Salvy, 2005)

(Chyzak, Mishna, 2026)

Deriving a Triple Sum (1/2)

We are interested in

$$\sum_{m,n=0}^{\infty} r_{m,n} q^m \frac{t^n}{n!} = \langle \exp(f(q, p)), \exp(g(p)t) \rangle$$

where

$$\langle p_1^{r_1} p_2^{r_2} p_3^{r_3}, p_1^{s_1} p_2^{s_2} p_3^{s_3} \rangle = \begin{cases} 1^{r_1} r_1! 2^{r_2} r_2! 3^{r_3} r_3! & \text{if } r = s, \\ 0 & \text{otherwise,} \end{cases}$$

$$f(q, p) = \frac{1}{2} q(p_1^2 - p_2) - \frac{1}{4} q^2 p_2^2 + \frac{1}{6} q^3 p_3^2, \quad g(p) = p_1 + \frac{1}{6} p_1^3 + \frac{1}{2} p_1 p_2 + \frac{1}{3} p_3.$$

Deriving a Triple Sum (1/2)

We are interested in

$$r_{m,n} = \left\langle [q^m] \exp(f(q, p)), g(p)^n \right\rangle$$

where

$$\langle p_1^{r_1} p_2^{r_2} p_3^{r_3}, p_1^{s_1} p_2^{s_2} p_3^{s_3} \rangle = \begin{cases} 1^{r_1} r_1! 2^{r_2} r_2! 3^{r_3} r_3! & \text{if } r = s, \\ 0 & \text{otherwise,} \end{cases}$$

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Deriving a Triple Sum (1/2)

We are interested in

$$r_{m,n} = ([q^m] \exp(f(q, p)) \odot g(1p_1, 2p_2, 3p_3)^n \odot K(p_1)K(p_2)K(p_3))_{\substack{p_1=1 \\ p_2=1 \\ p_3=1}}$$

where

$$p_1^{r_1} p_2^{r_2} p_3^{r_3} \odot p_1^{s_1} p_2^{s_2} p_3^{s_3} = \begin{cases} p_1^{r_1} p_2^{r_2} p_3^{r_3} & \text{if } r = s, \\ 0 & \text{otherwise,} \end{cases}$$

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$$K(x) = \sum_{\ell=0}^{\infty} \ell! x^\ell.$$

Deriving a Triple Sum (2/2)

We are interested in

$$r_{m,n} = \left([q^m] \exp(f(q, p)) \odot \tilde{g}(p)^n \odot K(p_1)K(p_2)K(p_3) \right)_{p_1=p_2=p_3=1}$$

where

$$f(q, p) = \frac{1}{2}q(p_1^2 - p_2) - \frac{1}{4}q^2p_2^2 + \frac{1}{6}q^3p_3^2, \quad \tilde{g}(p) = p_1 + \frac{1}{6}p_1^3 + p_1p_2 + p_3.$$

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By multiple applications of the binomial theorem:

$$\tilde{g}(p)^n = \sum_{a=0}^n \sum_{b=0}^{n-a} \sum_{c=0}^{n-a-b} \binom{n}{a, b, c, n-a-b-c} \frac{1}{2^b 3^b} p_1^{a+3b+c} p_2^c p_3^{n-a-b-c},$$

$$\begin{aligned} [q^m] \exp(f(q, p)) &= \sum_{u=0}^{\lfloor m/3 \rfloor} \sum_{v=0}^{\lfloor (m-3u)/2 \rfloor} \sum_{w=0}^{m-3u-2v} \frac{1}{u! v! w! (m-3u-2v-w)!} \\ &\quad \times \frac{(-1)^{m+u+v+w}}{2^{m-2u} 3^u} p_1^{2w} p_2^{m-3u-w} p_3^{2u}. \end{aligned}$$

Deriving a Triple Sum (2/2)

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$$a_k = r_{2k-1,2k} = \left([q^{2k-1}] \exp(f(q, p)) \odot \tilde{g}(p)^{2k} \odot K(p_1)K(p_2)K(p_3) \right)_{p_1=p_2=p_3=1}$$

where

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$$[q^m] \exp(f(q, p)) = \sum_{u=0}^{\lfloor m/3 \rfloor} \sum_{v=0}^{\lfloor (m-3u)/2 \rfloor} \sum_{w=0}^{m-3u-2v} \frac{1}{u! v! w! (m-3u-2v-w)!} \\ \times \frac{(-1)^{m+u+v+w}}{2^{m-2u} 3^u} p_1^{2w} p_2^{m-3u-w} p_3^{2u}.$$

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Forcing equal exponents in both sums yields

$$a_k = \sum_{u=0}^{\lfloor (2k-1)/3 \rfloor} \sum_{v=0}^{\lfloor (2k-1-3u)/2 \rfloor} \sum_{w=k-u}^{2k-1-3u-2v} \frac{(-1)^{u+v+w+1}}{2^{k-1-u+w} 3^{2u+w-k}} \\ \times \frac{(2k)! (2w)!}{(k+1)! (u+w-k)! u! v! w! (2k-1-3u-2v-w)!}.$$

Classical Creative Telescoping for Triple Sums

$$a_k = \sum_{u=0}^{\lfloor (2k-1)/3 \rfloor} \sum_{v=0}^{\lfloor (2k-1-3u)/2 \rfloor} \sum_{w=k-u}^{2k-1-3u-2v} f_{k,u,v,w}$$

What algorithms for creative telescoping compute and promise

(Chyzak, 2000) and (Koutschan, 2010) find a **telescoper** and **certificates**,

$$P \in \mathbb{Q}(k)\langle S_k \rangle, \quad Q_1, Q_2, Q_3, C_1, \dots, C_4 \in \mathbb{Q}(k, u, v, w)\langle S_k, S_u, S_v, S_w \rangle,$$

that satisfy

$$P = (S_u - 1)Q_1 + (S_v - 1)Q_2 + (S_w - 1)Q_3 + \sum_{\ell=1}^4 C_\ell Z_\ell$$

for annihilators of the summand,

$$Z_1 = S_k - \frac{f_{k+1,u,v,w}}{f_{k,u,v,w}}, \quad \dots, \quad Z_4 = S_w - \frac{f_{k,u,v,w+1}}{f_{k,u,v,w}}.$$

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that satisfy

$$\begin{aligned} P \cdot f_{k,u,v,w} &= (S_u - 1) \cdot (Q_1 \cdot f_{k,u,v,w}) + (S_v - 1) \cdot (Q_2 \cdot f_{k,u,v,w}) \\ &+ (S_w - 1) \cdot (Q_3 \cdot f_{k,u,v,w}) + \sum_{\ell=1}^4 C_\ell \cdot (Z_\ell \cdot f_{k,u,v,w}). \end{aligned}$$

Classical Creative Telescoping for Triple Sums

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Wishful motivation

$$\begin{aligned} P \cdot f_{k,u,v,w} &= (S_u - 1) \cdot (Q_1 \cdot f_{k,u,v,w}) + (S_v - 1) \cdot (Q_2 \cdot f_{k,u,v,w}) \\ &+ (S_w - 1) \cdot (Q_3 \cdot f_{k,u,v,w}) + \sum_{u,v,w} \sum_{\ell=1}^4 C_\ell \cdot (Z_\ell \cdot f_{k,u,v,w}) \end{aligned}$$

should imply, after summing over $\{0, \dots, N\}^3$ with large enough N ,

$$\begin{aligned} P \cdot \sum_{u,v,w=0}^N f_{k,u,v,w} &= \left(\sum_{v,w=0}^N Q_1 \cdot f_{k,u,v,w} \right)_{u=0}^{N+1} + \left(\sum_{u,w=0}^N Q_2 \cdot f_{k,u,v,w} \right)_{v=0}^{N+1} \\ &+ \left(\sum_{u,v=0}^N Q_3 \cdot f_{k,u,v,w} \right)_{w=0}^{N+1} + \sum_{u,v,w=0}^N 0 = 0. \end{aligned}$$

Classical Creative Telescoping for Triple Sums

$$a_k = \sum_{u=0}^{\lfloor (2k-1)/3 \rfloor} \sum_{v=0}^{\lfloor (2k-1-3u)/2 \rfloor} \sum_{w=k-u}^{2k-1-3u-2v} f_{k,u,v,w}$$

What really happens: many ill-shaped auxiliary sums \rightarrow intractable

$$\begin{aligned} P \cdot f_{k,u,v,w} &= (S_u - 1) \cdot (Q \cdot f_{k,u,v,w}) + (S_v - 1) \cdot (Q_2 \cdot f_{k,u,v,w}) \\ &+ (S_w - 1) \cdot (Q_3 \cdot f_{k,u,v,w}) + \sum_{u,v,w} \sum_{\ell=1}^4 C_\ell \cdot (Z_\ell \cdot f_{k,u,v,w}) \end{aligned}$$

- $Q, Q_2, Q_3, C_1, \dots, C_4, Z_1, \dots, Z_4$ have poles,
- absorbing poles by using shifts is **not a sufficient proof** (any syntactic change in equation requires an independent proof),
- Equation **need not hold everywhere**, even away from their poles (more poles may have been lost during the computer calculation),
- finite-difference operators **do not commute** with summation,
- summation range **has to be split** into sub-ranges.

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Similar observations in the past

- (Chyzak, Mahboubi, Sibut-Pinote, Tassi, 2014)
- (Koutschan, Wong, 2021)

Do not provide a sufficient solution.

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Do not provide a sufficient solution.

(Bostan, Lairez, Salvy, 2015) is not applicable

a_k is not a binomial sum (non-trivial denominator).

Towards Everywhere-Valid Recurrence Relations

Iverson's bracket notation

Definition:

$$\llbracket \mathcal{P} \rrbracket = \begin{cases} 1 & \text{if the predicate } \mathcal{P} \text{ holds,} \\ 0 & \text{if it does not hold.} \end{cases}$$

Nice property: if $c(e)$ is defined everywhere,

$$c(e) \llbracket e = 0 \rrbracket = c(0) \llbracket e = 0 \rrbracket.$$

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Everywhere-defined factorial and “inverse factorial” functions

$$n! = \begin{cases} n! & \text{if } n \geq 0, \\ 0 & \text{if } n < 0. \end{cases} \quad n_i = \begin{cases} 1/n! & \text{if } n \geq 0, \\ 0 & \text{if } n < 0. \end{cases}$$

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Iverson's bracket notation

Definition:

$$\llbracket \mathcal{P} \rrbracket = \begin{cases} 1 & \text{if the predicate } \mathcal{P} \text{ holds,} \\ 0 & \text{if it does not hold.} \end{cases}$$

Nice property: if $c(e)$ is defined everywhere,

$$c(e) \llbracket e = 0 \rrbracket = c(0) \llbracket e = 0 \rrbracket.$$

Everywhere-defined factorial and “inverse factorial” functions

$$n! = \begin{cases} n! & \text{if } n \geq 0, \\ 0 & \text{if } n < 0. \end{cases} \quad n_i = \begin{cases} 1/n! & \text{if } n \geq 0, \\ 0 & \text{if } n < 0. \end{cases}$$

Everywhere-valid recurrence relations

$$\forall n \in \mathbb{Z}, n! = n(n-1)! + \llbracket n = 0 \rrbracket$$

$$\forall n \in \mathbb{Z}, n_i = (n+1)(n+1)_i$$

Reformulation as a Sum over Natural Bounds

A triple sum with everywhere-defined summand

$$a_k = \sum_{u \in \mathbb{Z}} \sum_{v \in \mathbb{Z}} \sum_{w \in \mathbb{Z}} f_{k,u,v,w}$$

where

$$f_{k,u,v,w} = (-1)^{u+v+w+1} 2^{1+u-w-k} 3^{k-2u-w} \\ \times (2k)! (2w)! \times (k+1)_i (u+w-k)_i u_i v_i w_i (2k-1-3u-2v-w)_i.$$

The summand $f_{k,u,v,w}$ satisfies first-order recurrence relations valid over \mathbb{Z}^4 . But we won't try to use them directly.

Variant Creative Telescoping

Step 1: Heuristic creative telescoping

Obtain 4-variate rational functions Q_i such that

$$P(k, S_k) - \Delta_u Q_1(k, u, v, w) - \Delta_v Q_2(k, u, v, w) - \Delta_w Q_3(k, u, v, w) \equiv 0 \pmod{Z_1, \dots, Z_4}.$$

“Integer” denominator of, e.g., Q_1 : $(k+3) \prod_{i=2}^8 (2k+i-3u-2v-w)$.

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Step 2: Removal of denominators, by trial and error

Remove denominators by finding suitable shifts such that

$$P(k, S_k) - \Delta_u R_1(k, u, v, w) S_k^9 S_w^9 - \Delta_v R_2(k, u, v, w) S_k^{10} S_w^{10} - \Delta_w R_3(k, u, v, w) S_k^{10} S_w^{10} \equiv 0 \pmod{Z_1, \dots, Z_4}.$$

Common “integer” denominator of R_1, R_2, R_3 : $k+3$.

Variant Creative Telescoping

There remains to prove $\forall (k, u, v, w) \in \mathbb{Z}^4, k \neq -3 \implies z_{k,u,v,w} = 0$ for

$$\begin{aligned} z_{k,u,v,w} := & p_5(k)f_{k+5,u,v,w} + p_4(k)f_{k+4,u,v,w} + p_3(k)f_{k+3,u,v,w} \\ & + p_2(k)f_{k+2,u,v,w} + p_1(k)f_{k+1,u,v,w} + p_0(k)f_{k,u,v,w} \\ & - R_1(k, u+1, v, w)f_{k+9,u+1,v,w+9} + R_1(k, u, v, w)f_{k+9,u,v,w+9} \\ & - R_2(k, u, v+1, w)f_{k+10,u,v+1,w+10} + R_2(k, u, v, w)f_{k+10,u,v,w+10} \\ & - R_3(k, u, v, w+1)f_{k+10,u,v,w+11} + R_3(k, u, v, w)f_{k+10,u,v,w+10}. \end{aligned}$$

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Step 3: Rewrite each $f_{k+p,u+q,v+r,w+s}$ in terms of $f_{k+11,u,v,w+12}$ (not $f_{k,u,v,w}$!)

Determine and prove 12 everywhere-valid identities, like

$$\begin{aligned} \forall (k, u, v, w) \in \mathbb{Z}^4, \\ (k+11)(2k+21)(2w+23)f_{k+10,u,v,w+11} = \\ - (k+12)(2k+9-3u-2v-w)f_{k+11,u,v,w+12}. \end{aligned}$$

- Shift $f_{k+11,u,v,w+12}$ determined by trial and error.
- Introduces no new denominator in $z_{k,u,v,w}$.

Variant Creative Telescoping

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Proof postponed: **new algorithmic idea.**

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Step 4: Normalizing the candidate relation

$$z_{k,u,v,w} = \frac{1}{k+3} \frac{\text{polynomial expression}}{\text{polynomial expr. non-zero at integers}} \times f_{k+11,u,v,w+12} = 0$$

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Step 5: Telescoping the telescoping relation

For $k \in \mathbb{Z} \setminus \{-3\}$, summing $z_{k,u,v,w}$ for (u, v, w) over \mathbb{Z}^3 yields

$$p_5(k)a_{k+5} + p_4(k)a_{k+4} + p_3(k)a_{k+3} + p_2(k)a_{k+2} + p_1(k)a_{k+1} + p_0(k)a_k = 0.$$

Proving two-term relations between shifts of f

Goal

$$a(k, u, v, w)f_{k+p, u+q, v+r, w+s} - b(k, u, v, w)f_{k+11, u, v, w+12} = 0$$

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Simplifying by

$f_{k+p, u+q, v+r, w+s}$ = expression in factorials and inverse factorials,

$$n! = n(n-1)! + \llbracket n = 0 \rrbracket, \quad n_i = (n+1)(n+1)_i,$$

$$c(e) \llbracket e = 0 \rrbracket = c(0) \llbracket e = 0 \rrbracket$$

leads to a linear combination of terms of the form

$$t := \text{poly}(k, u, v, w) \times \prod_i \ell_i(k, u, v, w)! \times \prod_i \ell'_i(k, u, v, w)_i \times \llbracket e = 0 \rrbracket.$$

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Fact

$$t \neq 0 \Rightarrow \bigwedge_i (\ell_i(k, u, v, w) \geq 0) \wedge \bigwedge_i (\ell'_i(k, u, v, w) \geq 0)$$

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Fact

$$t = 0 \Leftrightarrow \bigwedge_i (\ell_i(k, u, v, w) \geq 0) \wedge \bigwedge_i (\ell'_i(k, u, v, w) \geq 0) \text{ unsatisfiable}$$

Proving two-term relations between shifts of f

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$$a(k, u, v, w)f_{k+p, u+q, v+r, w+s} - b(k, u, v, w)f_{k+11, u, v, w+12} = 0$$

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$$c(e) \llbracket e = 0 \rrbracket = c(0) \llbracket e = 0 \rrbracket,$$

then replacing with 0 those unsatisfiable expressions of the form

$$\text{poly}(k, u, v, w) \times \prod_i \ell_i(k, u, v, w)! \times \prod_i \ell'_i(k, u, v, w)_i \times \llbracket e = 0 \rrbracket,$$

proves the goal in all 12 cases of interest.

Proving unsatisfiability: algorithms exist

Satisfiability Modulo Theory + Linear Integer Arithmetic (SMT LIA)

Second Approach:
By a Residue Representation

Deriving a Residue Representation

We are interested in

$$\sum_{m,n=0}^{\infty} r_{m,n} q^m \frac{t^n}{n!} = \langle \exp(f(q, p)), \exp(g(p)t) \rangle$$

where

$$f(q, p) = \frac{1}{2}q(p_1^2 - p_2) - \frac{1}{4}q^2 p_2^2 + \frac{1}{6}q^3 p_3^2, \quad g(p) = p_1 + \frac{1}{6}p_1^3 + \frac{1}{2}p_1 p_2 + \frac{1}{3}p_3,$$

$$\langle p_1^{r_1} p_2^{r_2} p_3^{r_3}, p_1^{s_1} p_2^{s_2} p_3^{s_3} \rangle = \begin{cases} 1^{r_1} r_1! 2^{r_2} r_2! 3^{r_3} r_3! & \text{if } r = s, \\ 0 & \text{otherwise.} \end{cases}$$

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Formal Laplace transform, formal residue

$$\mathcal{L}(p_1^{r_1} p_2^{r_2} p_3^{r_3}) = \frac{r_1!}{p_1^{r_1+1}} \frac{r_2!}{p_2^{r_2+1}} \frac{r_3!}{p_3^{r_3+1}}, \quad \text{res} \left(\sum_{r \in \mathbb{Z}^3} c_r p^r \right) = c_{-1, \dots, -1}.$$

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Residue representation

$$\left\langle \exp(f), \exp(tg) \right\rangle = \text{res}_{p,q} \left(e^{f(p_1, p_2, p_3, q)} \mathcal{L}(e^{tg(1p_1, 2p_2, 3p_3)}) \right)$$

Deriving a Residue Representation

We are interested in

$$A(t) := \sum_{k=0}^{\infty} r_{2k-1,2k} \frac{t^{2k}}{(2k)!} = \text{even}_t \text{diag}_{q,t} \left\langle q \exp(f(q,p)), \exp(g(p)t) \right\rangle$$

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$$f(q,p) = \frac{1}{2}q(p_1^2 - p_2) - \frac{1}{4}q^2 p_2^2 + \frac{1}{6}q^3 p_3^2, \quad g(p) = p_1 + \frac{1}{6}p_1^3 + \frac{1}{2}p_1 p_2 + \frac{1}{3}p_3,$$

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Residue representation

$$A(t) = \text{even}_t \text{res}_{p,q} \left(e^{f(p_1, p_2, p_3, q)} \mathcal{L} \left(e^{q^{-1} t g(1p_1, 2p_2, 3p_3)} \right) \right)$$

Integration in D-Modules (Brochet, Chyzak, Lairez, 2026)

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Composition of (holonomy-preserving) operations

- $\exp(t\tilde{g}(p))$: a system of annihilators is

$$q\partial_{p_i} - t\frac{d\tilde{g}}{dp_i} \quad (1 \leq i \leq 3), \quad q^2\partial_q + t\tilde{g}, \quad q\partial_t - \tilde{g}$$

- $\mathcal{L}(\dots)$: change $p_i \rightarrow -\partial_{p_i}$ and $\partial_{p_i} \rightarrow p_i$, and think modulo $\ker(\text{res}_p)$
- $\exp(f(p)) \times \dots$: change $\partial_{p_i} \rightarrow \partial_{p_i} - \frac{df}{dp_i}$
- $\text{res}_{p,q}(\dots)$: compute the integral of a module

Integration in D-Modules (Brochet, Chyzak, Lairez, 2026)

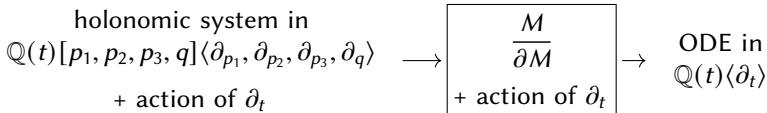
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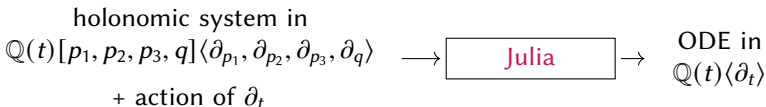
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Hadrien Brochet's `MultivariateCreativeTelescoping.jl`

Validity of the Algorithms

Role of holonomy

- Holonomic D-module $M \rightarrow$ its integral $M/\partial M$ is finite-dimensional.
- **Holonomic system** \rightarrow integration algorithm has **non-zero output**.

(Brochet, Chyzak, Lairez, 2026)

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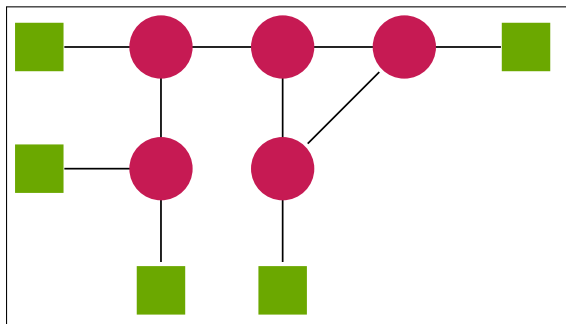
Creative telescoping for **D-finite functions is inappropriate** here

- D-finite functions \rightarrow work over $\mathbb{Q}(t, p_1, p_2, p_3, q)$.
- The system contains $p_3q - t$
 \rightarrow classical algorithms wrongly **believe they integrate 0**.

(Chyzak, 2000), (Koutschan, 2010)

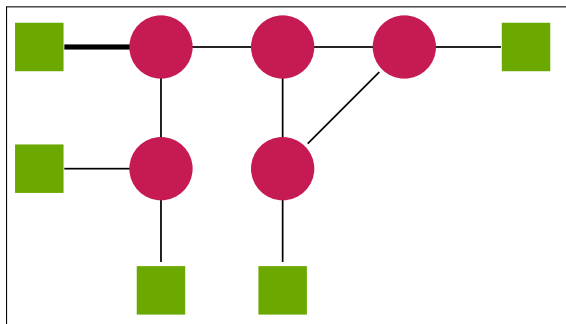
Third Approach:
By a Combinatorial Recurrence

Interrelated Combinatorial Classes



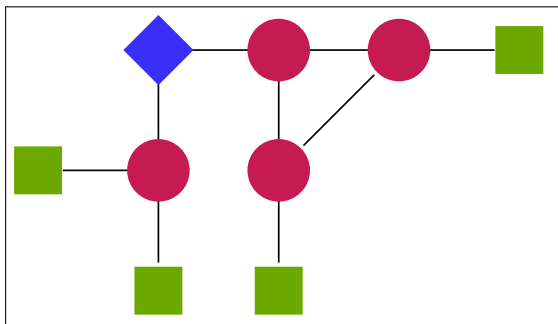
\mathcal{G}_0 | labelled graphs with vertices of degree 3 or 1 only

Interrelated Combinatorial Classes



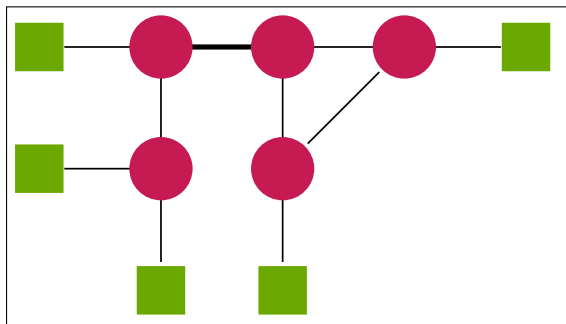
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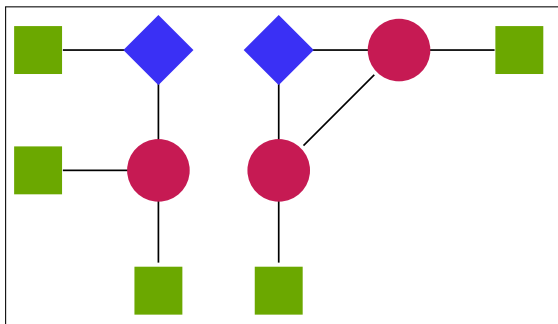
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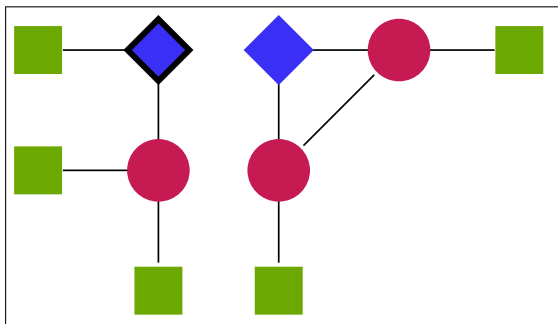
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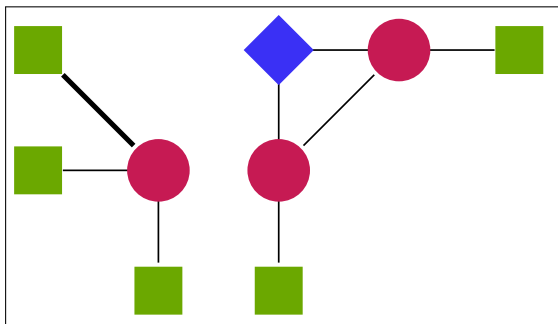
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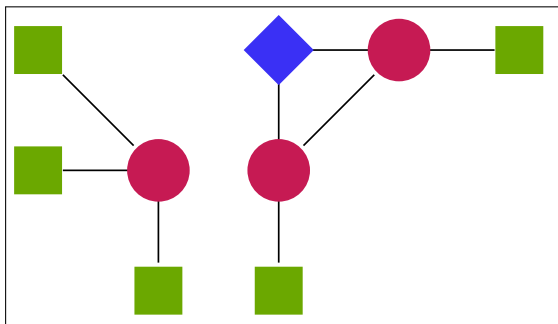
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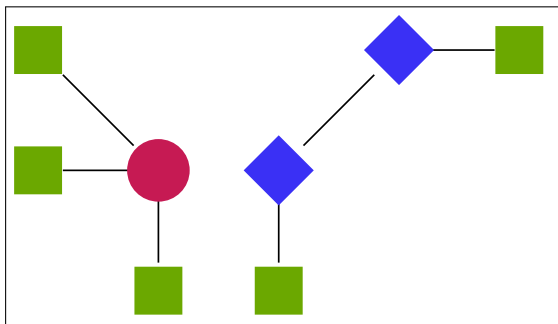
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Interrelated Combinatorial Classes



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A Differential System in Generating Functions

Four generating functions

$g_{a,b,c}$ = # of labelled graphs with a, b, c vertices of degree 1, 2, 3, $\tilde{g}_{a,2,c} = \dots$

$$G_b(q, t) = \sum_{a,c \geq 0} g_{a,b,c} q^{\frac{a+2b+3c}{2}} \frac{t^{a+b+c}}{(a+b+c)!}, \quad 0 \leq b \leq 2, \quad \tilde{G}_2(q, t) = \dots$$

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One of five decompositions

element of \mathcal{G}_0 with a marked edge \rightarrow element of \mathcal{G}_1 or element of $\tilde{\mathcal{G}}_2$

$$\begin{aligned} \frac{1}{2}(a+3c)g_{a,0,c} &= (a+c)g_{a-1,1,c-1} + \tilde{g}_{a,2,c-2} \\ q\partial_q \cdot G_0(q, t) &= qtG_1(q, t) + q\tilde{G}_2(q, t) \end{aligned}$$

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Similar differential equations for other classes or marking a vertex.

Algebraic Manipulations Get the Recurrence Relation

$$\partial_q \cdot G_0(q, t) = tG_1(q, t) + \tilde{G}_2(q, t),$$

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system for G_0 in $\mathbb{Q}(q, t)\langle\partial_q, \partial_t\rangle$

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↓ diagonal: creative telescoping for D-finite functions

differential equation for $\sum_{k \geq 0} a_k \frac{t^{2k}}{(2k)!}$

Conclusion

Creative telescoping in computational proofs

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traditionally: implicit assumption make **proofs incomplete**
here: fix whose **generality has to be explored**
- **residues:** CT in D-modules for residues
generally: ensuring a holonomic input system is **difficult**
here: polynomial exponential **make it simple**
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computer enumeration was used to **correct proof by cases**